Silicon IP Cores
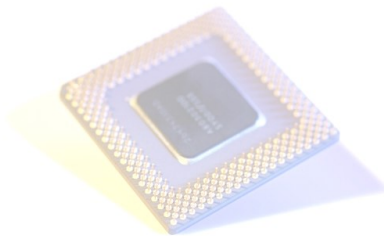
Hardware acceleration for seamless SoC integration

Product Catalog

May 2019

NOESIS
TECHNOLOGIES

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About Noesis Technologies P.C.

Noesis Technologies P.C. is a world wide leading silicon IP cores provider specialized in hardware implementation of high computational complexity telecom algorithms. Our hardware accelerator IP solutions allow telecom system developers to significantly off load demanding tasks from the CPU and to drastically decrease execution time thus boosting the overall system performance. Our IP cores present an industry leading combination of high performance, low power and low die-area, as well as easy customization for adaptability to a wide range of applications. Noesis offers a complete portfolio of Forward Error Correction IP core solutions that includes Reed Solomon Codecs, Viterbi Decoders, Turbo Product and Turbo Convolutional Codecs, LDCP Codecs, BCH codecs, (De)Interleavers, Channel Emulators. The company additionally offers a range of cores in the areas of security, networking, audio/voice/data compression, telecom DSP including a complete OFDM baseband processor.

Our company is also active in the development of integrated telecommunication systems that can used in education as well as in Research and Development applications. In the framework of this activity we have developed ComLab which is practically a cost efficient highly integrated development environment (IDE) that enables a system designer to rapidly build, configure and evaluate in real-time the performance of complex telecommunication systems. It is comprised of a Xilinx FPGA based board for the real-time HW emulation, a sophisticated application SW with interactive GUI capabilities for configuration, control and monitoring purposes as well as a rich portfolio of highly optimized telecom subsystem silicon IPs. ComLab platform is ideal for proof-of-concept rapid prototyping as well as an intuitive educational tool for engineers.

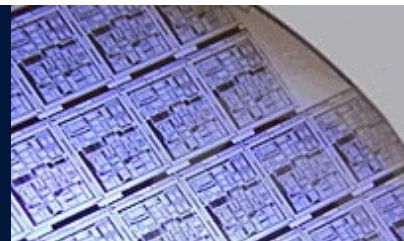Noesis Technologies also provides integrated solutions for WSN applications as well as develops disruptive technologies for the IoT market. In this framework, Noesis Technologies provides algorithms optimization and their effective implementation in FPGA technology using low power design techniques as well as embedded SW development for WSN nodes and application SW for base stations.

Noesis Technologies is a Xilinx Alliance Member and an official IP Core Designer for Turbo Coding technology. Its customer list includes U.S.A, Europe, Canada, Taiwan, China and India based companies active in telecom, defense and aerospace sectors. For further information please visit our web site www.noesis-tech.com

Forward Error Correction
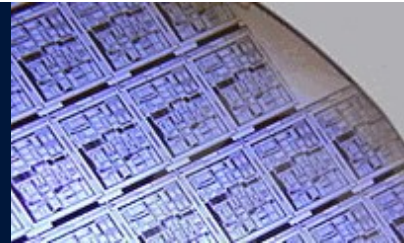
Error detection and correction codes are used nowadays in almost any digital transmission and storage system to ensure reliable transfer of information. Noesis Technologies offers a complete portfolio of forward error correction IP cores including Reed-Solomon, Viterbi, BCH, LDPC and Turbo Product codecs. These state-of-the-art hardware implementations have established Noesis Technologies as a worldwide leading provider in FEC IP core solutions. Noesis Technologies forward error correction IPs are developed to boost performance in wireless LANs and Internet, satellite communications, wireless broadband systems, optical networking, wireline access networks, data storage as well as in a variety of other target applications that require error detection and correction techniques. Noesis FEC solutions are highly configurable, with optimized architectures that can be targeted to multiple wireless or wired standards and can meet the most demanding area and speed application requirements. All of the FEC IP cores are silicon proven (FPGA or ASIC) and technology independent for easy and risk-free porting to any target silicon process.

Reed Solomon Codes

One of the most powerful and widespread used algebraic error correction block code is Reed Solomon algorithm. It belongs to the family of maximum distance separable non-binary linear cyclic codes and it is exceptionally powerful when channels errors occurs in bursts. It is also good at correcting random errors.

Noesis Technologies Reed Solomon highly parameterized hardware IP core solution can be used in any application that requires an RS based error correction control system. Its outstanding parameterization capability includes features such as configurable number of bits per symbol, maximum codeword length and maximum number of parity symbols. It also supports varying on the fly shortened and punctured codes. As a result of this parameterization and programmability feature any desirable code-rate can be easily achieved rendering the codec ideal for fully adaptive FEC applications. Erasures decoding for doubling the error correction capability as well as extended RS codes, burst or continuous decoding are also supported. The Noesis Technologies ntRS core has been specifically designed with fine-grain granularity in order to allow data path slices reconfigurability. The end-user can configure the number of slices in the data-path array in order to achieve the optimum ratio of throughput-rate vs silicon area and thus resulting into highly efficient hardware implementations.

Features
Fully configurable, time-domain, high throughput, Reed Solomon Decoder.
Supports different Reed Solomon coding standards.
Variable on the fly code rate adaptation by varying codeword length and/or number of parity symbols.
Variable bits per symbol, odd or even number of parity symbols.
Variable codeword length on a codeword by codeword basis.
Variable number of errors corrected on a codeword by codeword basis.
Supports shortened, punctured and extended codes.
Parameterized architecture allows optimum ratio of throughput-rate vs silicon area.
User configured primitive polynomial.
User configured generator polynomial.
Single or multiple symbol rate clock.
Continuous decoding with no gaps between codewords.
Predictable latency.
Counts number of errors and flags uncorrectable codewords.
Fully synchronous design, using single clock.
Silicon proven in ASIC and FPGA technologies for a variety of applications.

Forward Error Correction

ntRSE

Fully Configurable Reed Solomon Encoder
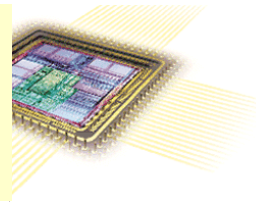
The ntRSE core implements the Reed Solomon encoding algorithm and is parameterized in terms of bits per symbol, maximum codeword length and maximum number of parity symbols. It also supports varying on the fly shortened codes. Therefore any desirable code-rate can be easily achieved rendering the encoder ideal for fully adaptive FEC applications. The ntRSE core supports continuous or burst encoding. The implementation is very low latency, high speed with a simple interface for easy integration in SoC applications.
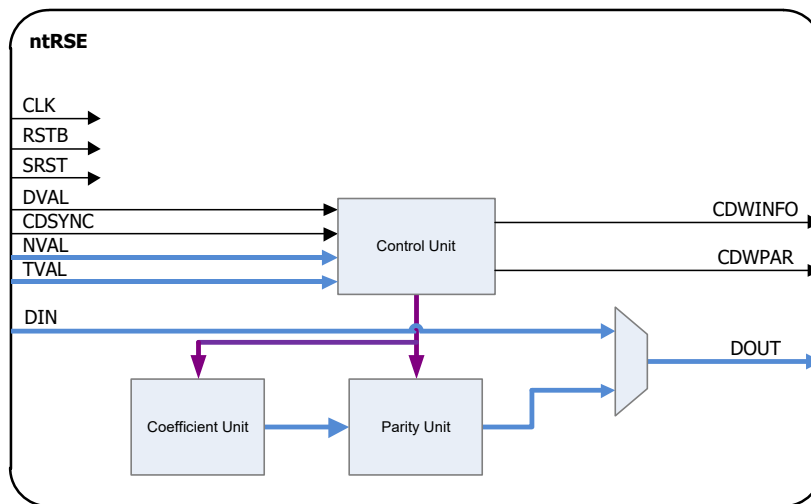
The ntRSE core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Virtex-2	RS(255,239)	585 CLB Slices	167
Altera	Stratix-II	RS(255,239)	333 ALUTs	162
TSMC	0.18 µm	RS(255,239)	2500 gates	250

ntRSD
Fully Configurable Reed Solomon Decoder

The ntRSD core implements a time-domain Reed-Solomon decoding algorithm. The core is parameterized in terms of bits per symbol, maximum codeword length and maximum number of parity symbols. It also supports varying on the fly shortened codes. Therefore any desirable code-rate can be easily achieved rendering the decoder ideal for fully adaptive FEC applications. The ntRSD core supports erasure decoding thus doubling its error correction capability. The core also supports continuous or burst decoding. The implementation is very low latency, high speed with a simple interface for easy integration in SoC applications.
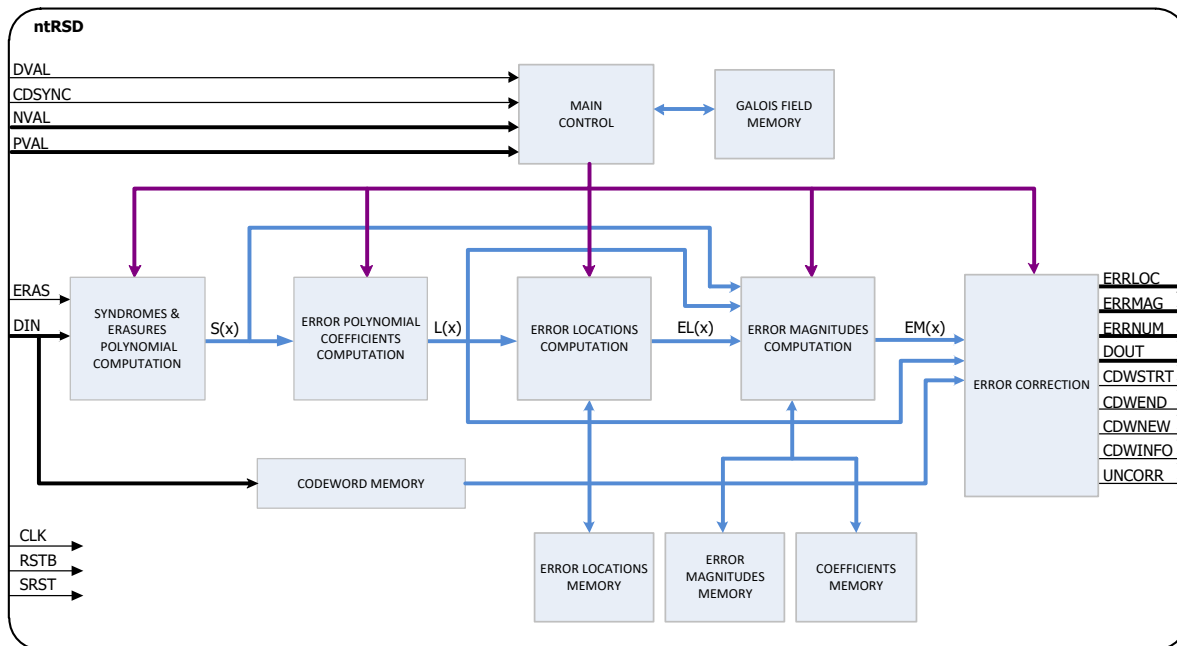
The ntRSD core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Virtex-2	RS(255,239)	2765 CLB Slices / 3 Block RAMs	88
Xilinx	Virtex-5	RS(255,239)	1490 CLB Slices / 3 Block RAMs	174
Xilinx	Spartan-3	RS(255,239)	2810 CLB Slices / 3 Block RAMs	50
Altera	Stratix-GX	RS(255,239)	5865 LCs / 3 Block RAMs	83
TSMC	0.18 μm	RS(255,239)	25 K gates / 12 K RAM bits	200

Forward Error Correction

Convolutional Codes
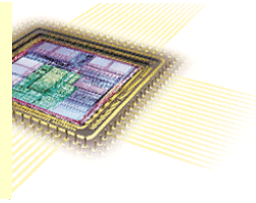
Linear convolutional codes are very well suited for random channel errors correction. When combined with Viterbi decoding algorithm can exploit the soft decision information provided by the demodulator and thus taking advantage of the 2dB coding again when compared with hard decision decoding. In application cases where noise is predominantly Gaussian then when the convolutional codes are concatenated with block codes result in an extremely powerful error correction control system.

Features
Fully configurable, high throughput convolutional FEC system based on Viterbi Decoder algorithm.
Supports different convolutional coding standards.
Parameterizable constraint length, code rate, generator coefficients and soft bits.
Parameterizable puncturing for full code rate control.
Programmable traceback depth.
Supports zero terminating and tail biting Viterbi decoding algorithm.
Soft or hard decision decoding.
Supports both continuous and burst input data flow.
Supports both block and continuous based decoding.
Fixed Viterbi decoder latency.
Single or multiple symbol rate clock.
Continuous decoding with no gaps between codewords.
Predictable decoder latency.
Area efficient design.
Fully synchronous design, using single clock.
Silicon proven in ASIC and FPGA technologies for a variety of applications.

ntVIT
Fully Configurable Viterbi Decoder
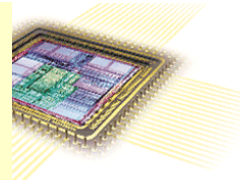
Convolutional FEC codes are very popular because of their powerful error correction capability and are especially suited for correcting random errors. The most effective decoding method for these codes is the soft decision Viterbi algorithm. ntVIT core is a high performance, fully configurable convolutional FEC core, comprised of a 1/N convolutional encoder, a variable code rate puncturer/depuncturer and a soft input Viterbi decoder. Depending on the application, the core can be configured for specific code parameters requirements. The highly configurable architecture makes it ideal for a wide range of applications. The convolutional encoder maps 1 input bit to N encoded bits, to generate a rate 1/N encoded bitstream. A puncturer can be optionally used to derive higher code rates from the 1/N mother code rate. On the encoder side, the puncturer deletes certain number of bits in the encoded data stream according to a user defined puncturing pattern which indicates the deleting bit positions. On the decoder side, the depuncturer inserts a-priori-known data at the positions and flags to the Viterbi decoder these bits positions as erasures. The Viterbi decoder uses a maximum-likelihood detection recursive process to correct errors in the data stream. The Viterbi input data stream can be composed of hard or soft bits. Soft decision achieves a 2 to 3dB increase in coding gain over hard-decision decoding. Data can be received continuously or with gaps.
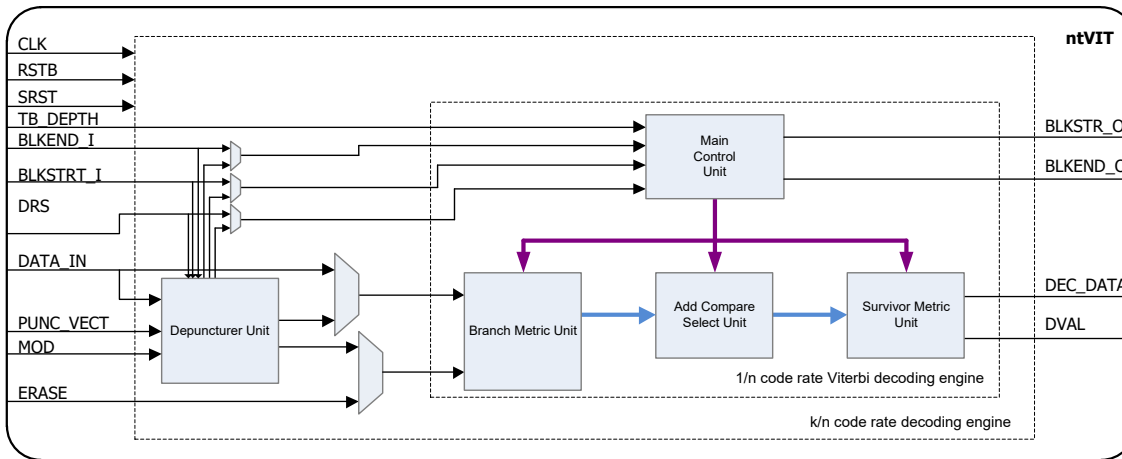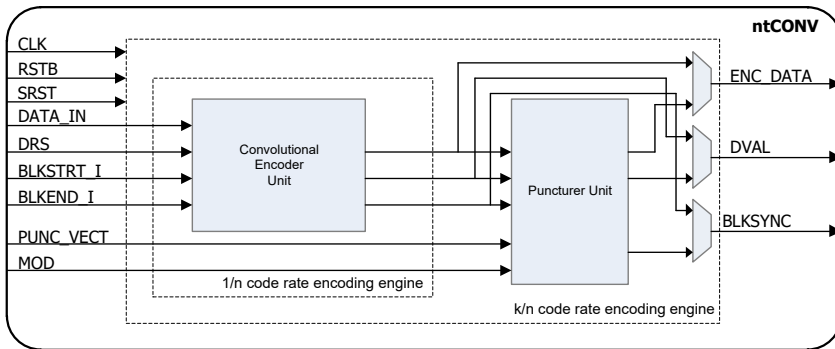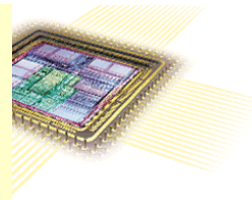

The ntVIT core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Virtex-5	1/2 rate, constraint length 7	2200 CLB Slices / 4 Block RAMs	150
Altera	Stratix-III	1/2 rate, constraint length 7	7384 ALUTs / 8 M144K Block RAMs	100
TSMC	180 nm	1/2 rate, constraint length 7	50K gates / 9K RAM bits	230

Turbo Product Codes
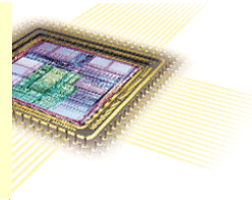
Turbo Product Codes (TPCs) exhibit excellent performance in moderate to high signal to noise ratios. Since TPCs have more advantage when a high rate code is used, they are ideal for commercial applications in wireless and satellite communications. Noesis Technologies ntTPC Turbo Product Codec solution is consisted of the Turbo Product Encoder (ntTPCE) and the Turbo Product Decoder (ntTPCD) IP cores. The ntTPC cores can be used in a variety of applications, including, wireless broadband communications, optical transmission systems, free space optical communication,
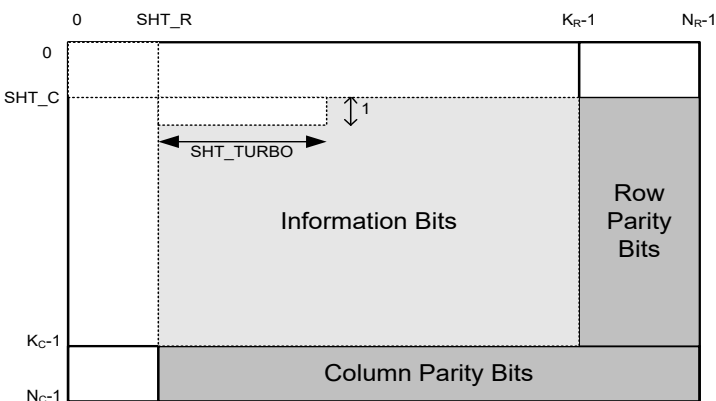
satellite modems.

The product code C is derived from two/three constituent codes, namely C1, C2 and optionally C3, thus supporting 2D or 3D codes. The information data is encoded in two/three dimensions. Every row of C is a code of C2 and every column of C is a code of C1. When the third coding dimension is enabled, then there are C3 C1*C2 data planes. The ntTPC cores support both e-Hamming and Single Parity codes as the constituent codes. The cores also support shortening of rows or columns of the product table, as well as turbo shortening. Shortening is a way of providing more powerful codes by removing information bits from the code. The construction of a two dimensional (N_{C-SHT_C}, K_{C-SHT_C}) × (N_{R-SHT_R}, K_{R-SHT_R}) code derived from the original (N_C, K_C) × (N_R, K_R) code is shown in the figure.

Features

Features
Encoder and decoder, support extended Hamming (256,247), (128,120), (64,57), (32,26), (16,11), (8,4) and Single Parity (256,255), (128,127), (64,63), (32,31), (16,15), (8,7) constituent error correcting codes.
3D encoding/decoding support with Single Parity (4,3) constituent code.
Highly programmable and parameterizable cores in terms of error correction capability, code rate, decoding iterations, decoding test patterns and scalability of design architecture.
Minimum system resources utilization and maximum resources reuse with one Single TPC elementary decoder instance for low power applications, producing up to 10Mbps information throughput (100Mhz − 4 decoding iterations).
Small area footprint of the elementary decoder also allows an alternative high throughput design approach with a number of cascaded / parallelized elementary decoders (plus the extra memories overhead), in order to reduce internal data re-iterations.
Flexible generic architecture with various combinations of parallelism options providing any desired application trade-off between area, performance and throughput rates.
Decoding algorithm achieves competitive performance results with the minimum possible test patterns and decoding iterations.
Bit serial encoder input/output interface. Soft input – soft output (SISO) serial decoder interface.
Flexible and programmable code rates, ranging from 0.1875 to 0.9922 (without shortening).
Additional programmability support for shortening of any selected code rate.
Programmable number of algorithmic iterations.
Simple yet robust encoder and decoder cores interface for optimum data flow control.
Synchronous single clock design.
Silicon proven in ASIC and Xilinx FPGA implementation technologies.

ntTPCE

Configurable Turbo Product Code Encoder
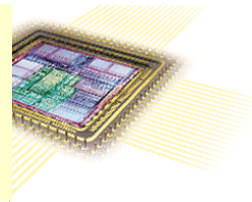
The ntTPCE core receives the information bits row by row from left to right and transmits the encoded bits in the same order. It consists of a row, column and 3D encoder. The row encoder encodes the data row-wise (C2). The encoded data produced from the row encoder are stored in an intermediate memory and reordered in a column-wise fashion. Once a full column has been written in the memory, the data are encoded column-wise by the column encoder (C1). When 3d encoding is employed, the encoded data produced from the column encoder are stored in an intermediate memory and reordered in a 3d-plane-wise fashion. The C3 data planes are encoded by an SPC(4,3) encoder (C3). Before output encoded information data are being reordered in row-wise fashion.
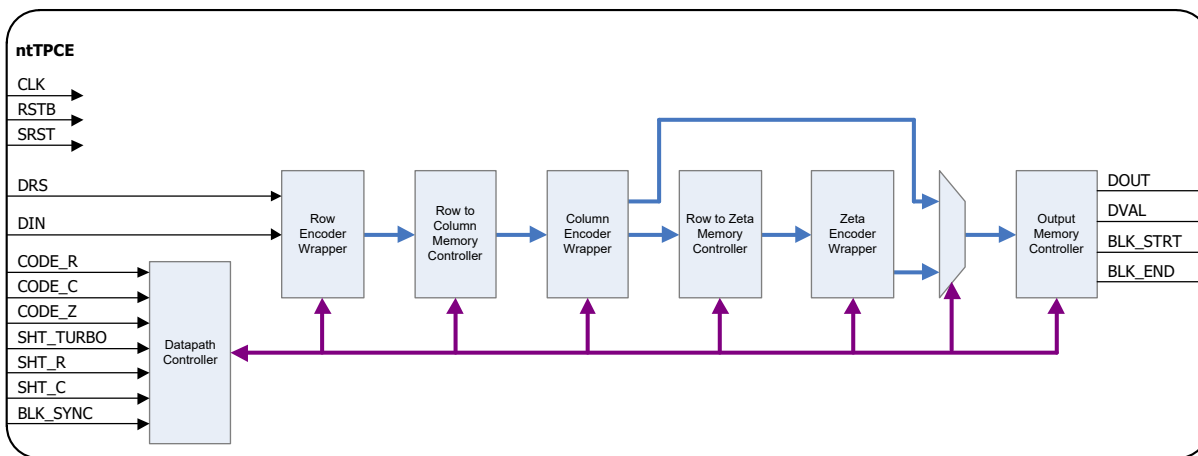
The ntTPCE core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details for the configurable ntTPCE core, as shown in the table range from the 64 bits/word—2D encoder to 256 bits/word—3D encoder configurations.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	64 bits/word—2D	334 Slices / 9216 Memory Bits	224
Xilinx	Kintex 7 XC7K410T-2	256 bits/word—3D	439 CLB Slices / 215040 Memory Bits	208

ntTPCD
Configurable Turbo Product Code Decoder
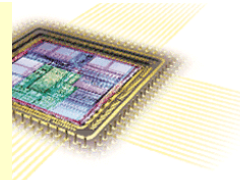
The ntTPCD decoder receives soft information from the channel in the 2's complement number system and the input samples are received row by row from left to right. The decoded soft information is output in the same order. The implemented decoding algorithm computes the extrinsic information for every dimension C1, C2, C3 by iteratively decoding words that are near the soft-input word. These words are called test patterns and their number is pre-configurable. All C1, C2, C3 words decoding takes place in a main decoding unit, the programmable elementary Soft Input Soft Output (SISO) decoder. An advanced scalable and parametric design approach produces custom design versions tailored to end customer applications design tradeoffs.

The architecture of one elementary SISO decoder shown below is parameterizable in terms of maximum constituent code size (64,128,256 bits), optional 3D codes support and maximum parallel test patterns processing (8,16,32) and soft bits.
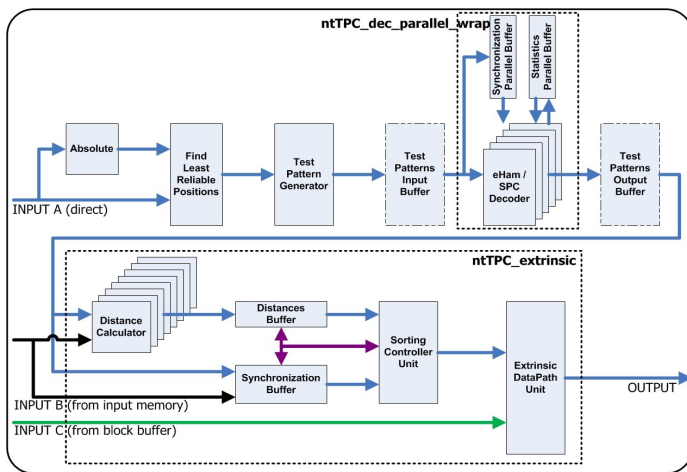
The following figures demonstrate kintex-7 device resources utilization for various configurations of the single SISO decoder architecture.
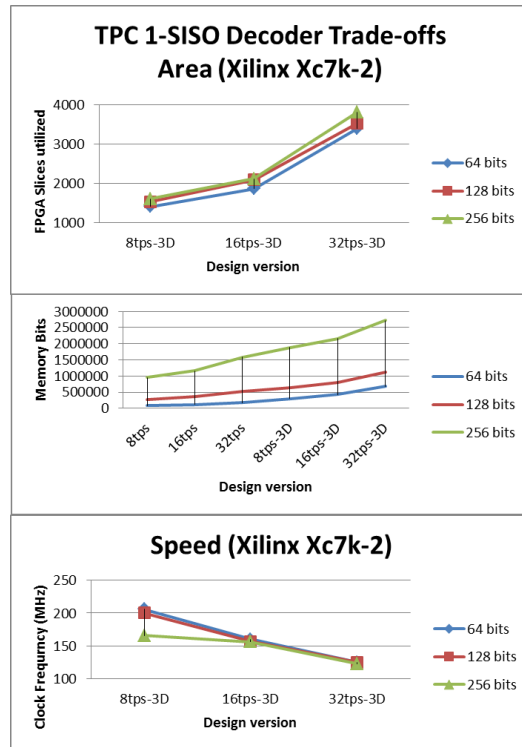
Depending on system trade-offs / requirements, one or more SISO decoders may be used in one of the following schemes:
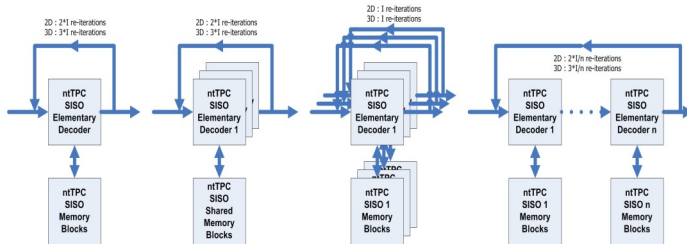
The BER vs SNR performance of the ntTPCD for various code rates, QPSK modulation, 8/16/32 test patterns and 2D/3D decoding are demonstrated in the following curves:

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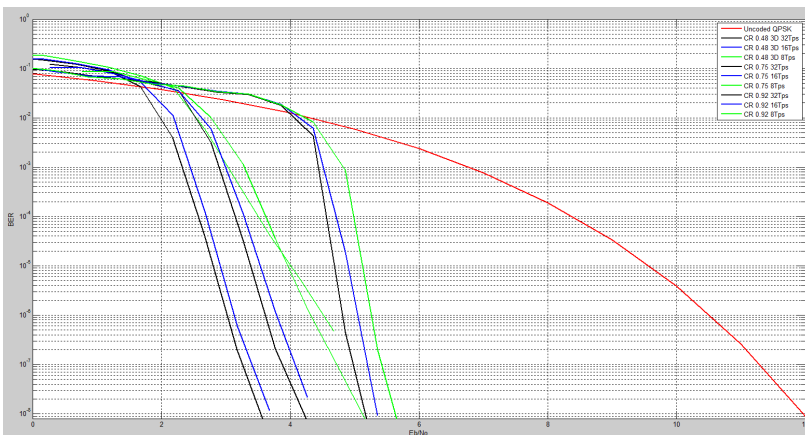
ntCTCE
HomePlug AV2 CTC Encoder
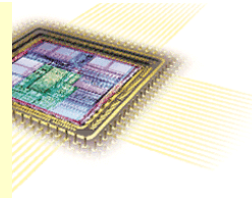
The ntCTCE encoder core uses two RCS constituent encoders, an Interleaver and an optional Puncturer to encode the u1 and u2 systematic input bits and to produce the x1 and x2 parity bits. When puncturing is not used all u1, u2, x1 and x2 are of the same size L and the overall code rate is 1/2. When puncturing is used, then x1 and x2 are punctured and the overall code rate is 16/18. Input and output of both ntCTCE and ntCTCD appear in natural order and the specification required bit order modifications are performed internally in each IP core. Information is partitioned in packets of data, the sizing of which is defined by the active mode of operation. The natural order of data within a packet is defined incrementally from bit 0 (b0). The ntCTCE DIN input port requires 2 bits in parallel, where the LSB should be b0, b2, b4, … and the MSB should be b1, b3, b5, ….

After the input data up for transmission are being scrambled, the scrambled data enter the ntCTCE encoder and parity is generated. The encoded data enter the Channel Interleaver and are output concatenated in groups of 4 bits, partly due to the nature of the Interleaving algorithm, and partly due to the necessity to maintain the TX system data rate.
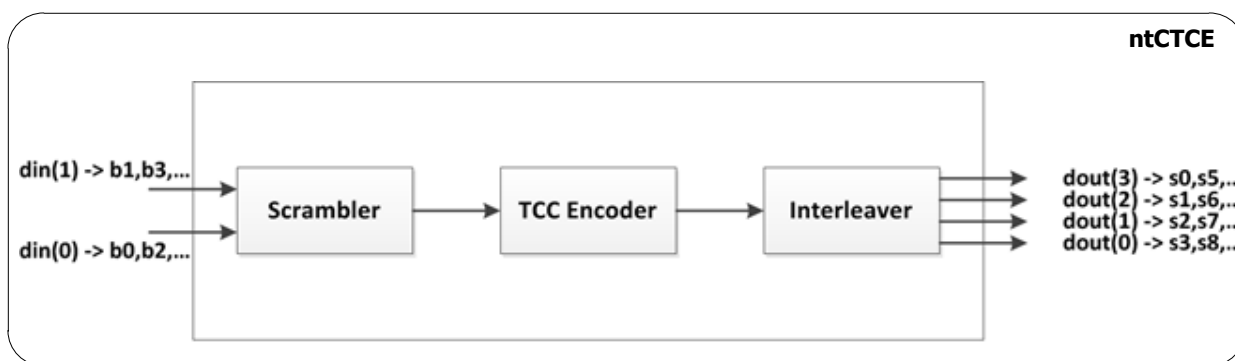
The following table demonstrates the mode of operation and relative sizing information.

Mode	Code Rate	Size L	Encoded bits	U1, U2 bits	X1, X2 bits
0	1/2	64	256	64	64
1	1/2	288	1152	288	288
2	1/2	544	2176	544	544
3	1/2	1056	4224	1056	1056
4	16/18	544	1224	544	68
5	1/2	2080	8320	2080	2080
6	RESERVED	-	-	-	-
7	16/18	2080	4680	2080	260

The ntCTCE core has been synthesized using Xilinx ISE Design Suite tools and Altera Quartus tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device and Aria V GX 5AGXFB3H4F35C5 device with a default balanced optimization strategy between area and timing. The area and performance metrics produced are summarized in the following tables.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	HomePlug AV2	198 Slices / 4 Block RAMS	201
Xilinx	Aria V GX 5AGXFB3H4F35C5	HomePlug AV2	1954 ALMs / 24704 Memory Bits	258

Forward Error Correction

ntCTCD
HomePlug AV2 CTC Decoder
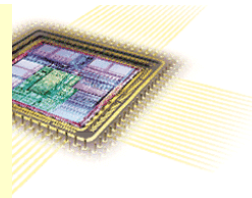
The ntCTCD decoder implements a Depuncturer, two MAP decoders, an Interleaver, a Deinterleaver and a hard decision unit. Each MAP decoder calculates log domain extrinsic probabilities and passes them to the next MAP decoder with the necessary interleaving of de-interleaving transformations. This procedure is repeated for a number of decoding iterations and each time the TCC decoder improves its estimate on the received bits. At the end of the predefined number of decoding iterations the decoder performs a hard decision and outputs the decoded bits. The received channel samples are scaled channel LLRs or quantized SOFT values (signed S8.3 fixed point format). The ntCTCD IP core needs these samples to be concatenated and ordered in groups of 4, in the same way as the ntCTCE output. Input samples are provided to the Deinterleaver, which in turn modifies the data stream to feed the one or multiple decoder parallel processor instances. Therefore the user can achieve the application target throughput rate by selecting the appropriate number of concurrently operating decoder instances. Once iterative turbo decoding has been performed, the hard decision resulting bits are de-scrambled and returned to IP core output in the same natural order, as they had been provided to ntCTCE input.
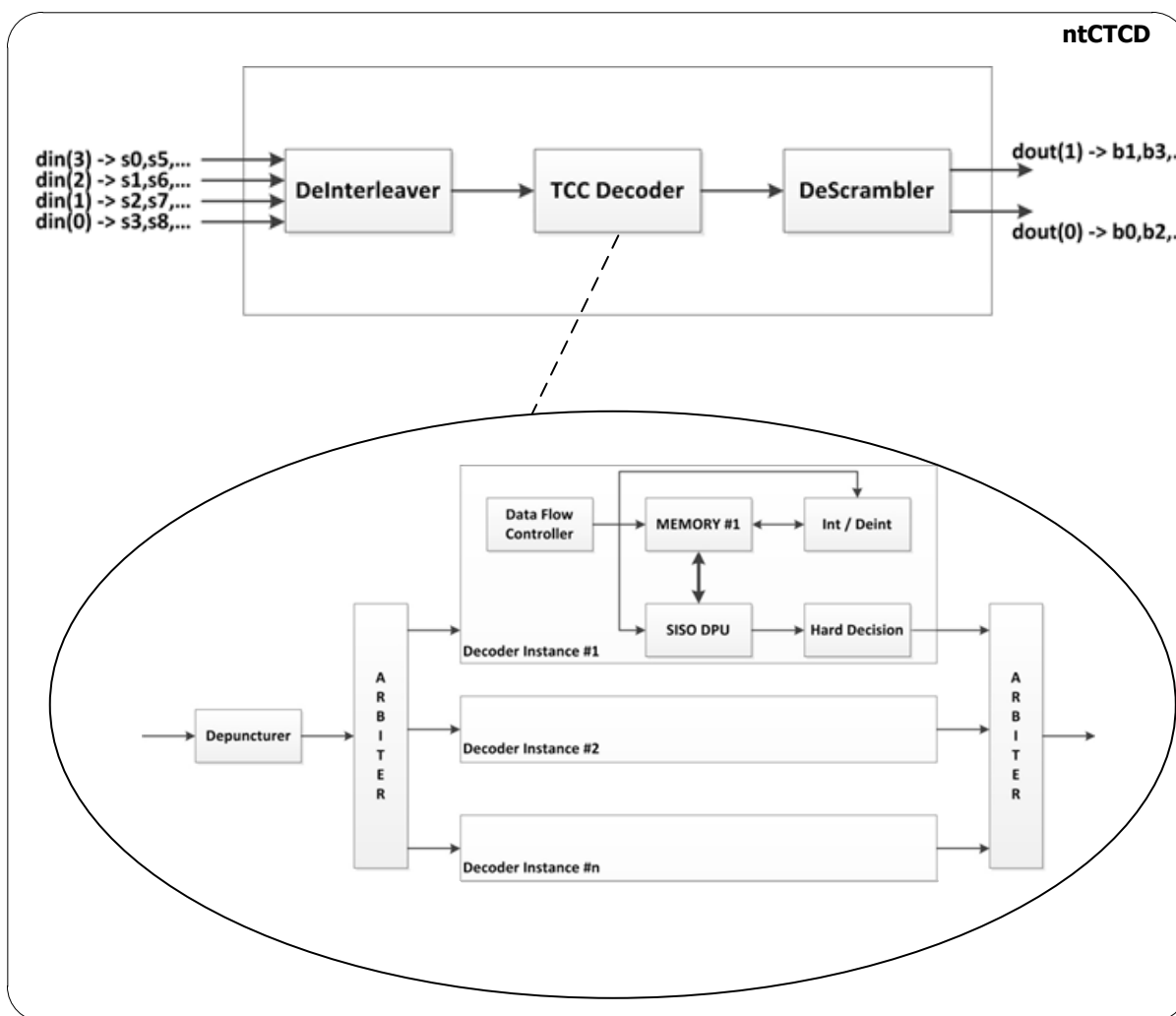
The ntCTCD core has been synthesized using Xilinx ISE Design Suite tools and Altera Quartus tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device and Aria V GX 5AGXFB3H4F35C5 device with a default balanced optimization strategy between area and timing. The area and performance metrics produced are sum-marized in the following tables.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	HomePlug AV2	2330 Slices / 41 Block RAMS	80
Xilinx	Aria V GX 5AGXFB3H4F35C5	HomePlug AV2	5708 ALMs / 798848 Memory Bits	72

ntCTCD
HomePlug AV2 CTC Decoder
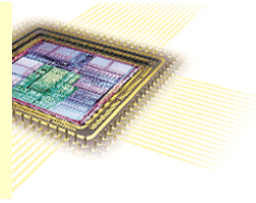
The ntCTCD achieves exceptional error correction performance as illustrated in the following BER vs Es/No graphs for the 7 modes of operation as described in HomePlug AV2 specification. The measurement conditions were as follows:

□ Input LLR S8.3 (signed, 5 integer bits, 3 fractional bits)

□ AWGN channel impairments

□ QPSK modulation

□ Approximate LLR calculation with proportional σ**2 scaling

□ 10 decoding iterations

QPSK BER vs Es/N0 - 1000 errors - 10 iterations - S8.3 Input LLR
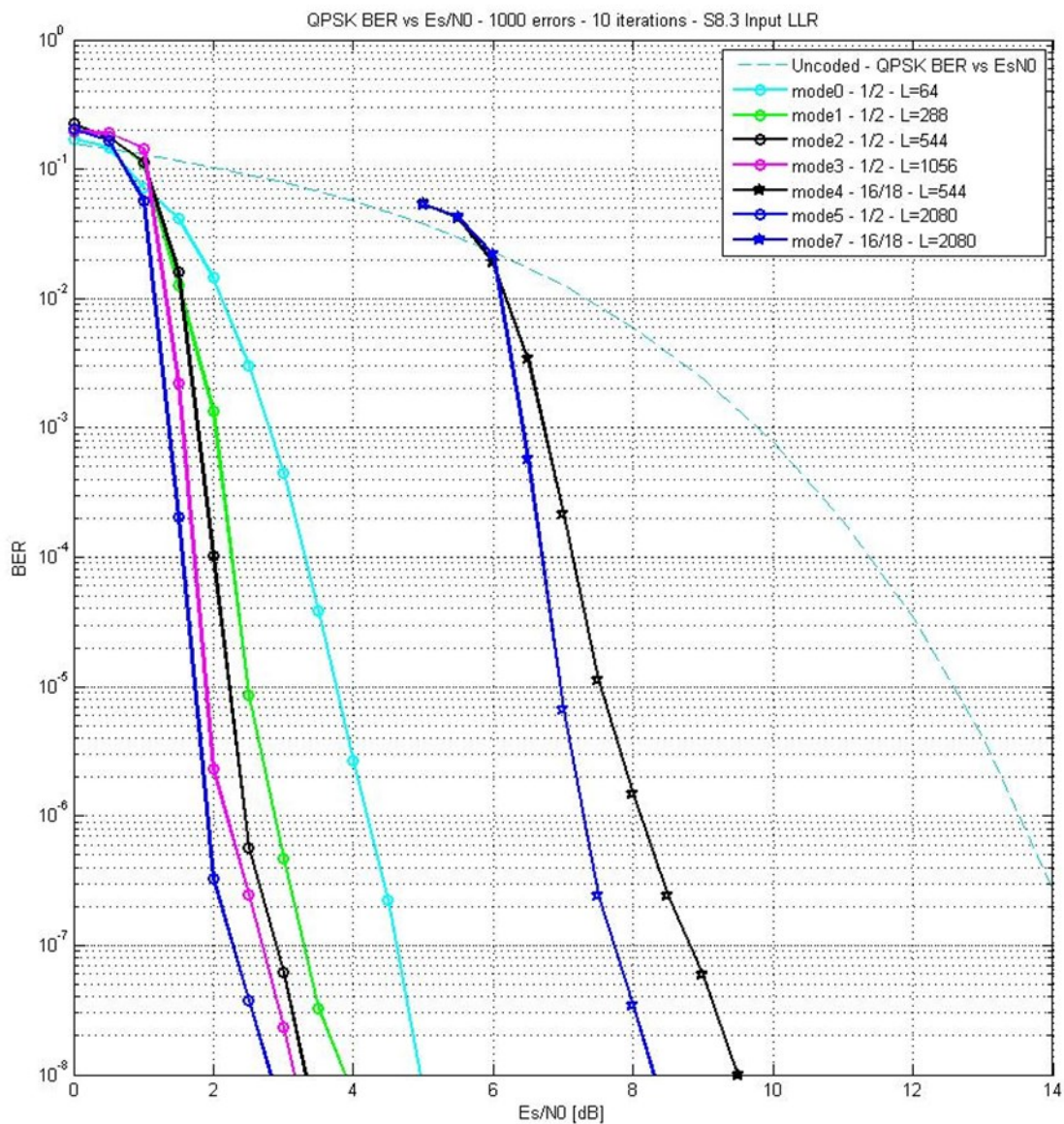
Low Density Parity Check Codes
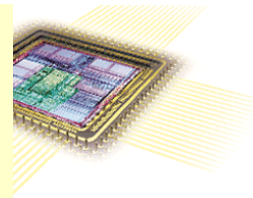
Low-density parity-check (LDPC) codes were introduced in 1960 at MIT by Robert G. Gallager in his Phd Dissertation. Low-density parity-check (LDPC) codes are a class of linear block codes. The name comes from the characteristic of their sparse parity-check matrix. The decoding of LDPC codes is done through an iterative, information message passing process. However, due to their computational effort and implementation complexity they were largely ignored till 1995 when D. McKay and R. Neal "rediscovered" them. Since then, many modern telecommunications systems have adopted LDPC codes as their coding scheme. The LDPC codes have excellent performance, which is very close to the channel capacity limit as defined by Claude E. Shannon Theorem. They offer reliable data transmission, particularly in noisy telecommunications channels. For their hardware implementation, the high parallelism degree they offer plays an important role.

Noesis Technologies has designed a highly efficient, modular and patent pending VLSI architecture of a certain type of structured LDPC Codes called Quasi-Cyclic LDPC Codes (QC-LDPC) or LDPC Block Codes (LDPC-BC). These LDPC codes are suitable for efficient hardware implementation and are based on block-structured LDPC codes with circular block matrices. The parity check matrix designed in this way can be conveniently represented by a base (block) matrix. This form of the parity check matrix simplifies the encoding and the decoding procedure. As a result, the main advantage is that they offer high throughput at low implementation complexity and they are considered in many applications and communication standards. The ntLDPCE (encoder) and ntLDPCD (decoder) cores can be used in a variety of applications, including but not limited to:

- Next generation Wired Home: Networking G.9960/G.9961 (G.hn).
- Digital Video Broadcasting: DVB-S2, DVB-S2X, DVB-T2, DVB-C2.
- Deep-space satellite missions (CCSDS).
- WiMax (IEEE 802.16e).
- WiFi (IEEE 802.11n - IEEE 802.11ac).
- WiGig (IEEE 802.11ad).
- WPAN (IEEE 802.15.3c).
- Hard disks.
- 10 Gigabit Ethernet - 10GBASE-T (IEEE 802.3an).
- CMMB (China Multimedia Mobile Broadcasting).

Features
Near Shannon limit performance.
Patent pending, highly efficient and modular hardware implementation.
Simple encoding and decoding procedure due to adoption of LDPC-BC.
Support of variable sub-matrix sizes (Z) of LDPC-BC.
Expandable parallelism degree based on the sub-matrix sizes (Z).
The layer scheduling of the decoding algorithm in tandem with the stopping decoding criterion offer fast convergence.
Fully configurable, high throughput, low cost implementation.
Ability to support different communication standards and a variety of practical applications with minor architectural modifications.
High flexibility in terms of code rates, decoding iterations, data width.
Adjustable trade-off between performance, throughput and area.
Flexible interface for easy system integration.
Fully synchronous design, using single clock.
Silicon proven in ASIC and Xilinx FPGA implementation technologies.

ntLDPCE-Ghn

G.hn Low Density Parity Check Encoder

The ntLDPCE-Ghn core implements the Quasi-Cyclic LDPC Block Codes (QC-LDPC-BC). These LDPC codes are based on block-structured LDPC codes with circular block matrices. The entire parity check matrix can be partitioned into an array of block matrices, each block matrix is either a zero matrix or a right cyclic shift of an identity matrix. The parity check matrix designed in this way can be conveniently represented by a base (block) matrix. The main advantage is that they offer high throughput at low implementation complexity and they are used in many applications and communication standards. The ntLDPCE-Ghn core is fully compliant with various wireless and wireline communication standards including ITU-T G.9960 (G.hn), IEEE 802.16e (WiMAX), IEEE 802.11n/ac (WiFi) etc. The core is highly reconfigurable and is able to support different sub-matrix sizes (Z) of LDPC-BC, that are tailored for specific applications. It also supports varying on the fly code rates and input data width. The implementation is flexible, high speed, efficient area utilization and has a simple interface for easy integration in SoC applications.
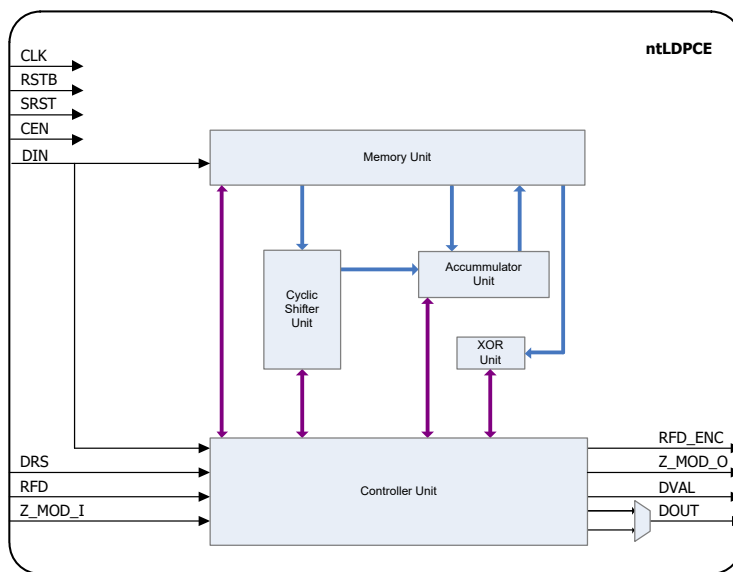
The ntLDPCE-Ghn core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details and performance metrics of the ntLDPCE core configured for ITU-T 9960 G.hn standard are shown in the tables below.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	ITU-T 9960 G.hn compliant	4250 Slices / 31 Block RAMs	203

Mode	Code Rate	(code_len, info_len)	Throughput Rate (Gbps) Z-bits parallelism	Throughput Rate (Gbps) 8-bits parallelism
1	$(1/2)_H$	(336,168)	0.310	0.272
2	$(1/2)_S$	(1920,960)	1.355	0.748
3	$(1/2)_L$	(8640,4320)	1.458	0.745
4	$(2/3)_S$	(1439,960)	1.397	0.998
5	$(2/3)_L$	(6480,4320)	1.447	0.980
6	$(5/6)_S$	(1152,960)	1.442	1.249
7	$(5/6)_L$	(5184,4320)	1.479	1.243

ntLDPCD-Ghn

G.hn Low Density Parity Check Decoder

The ntLDPCD-Ghn core implements the Quasi-Cyclic LDPC Block Codes (QC-LDPC-BC). These LDPC codes are based on block-structured LDPC codes with circular block matrices. The entire parity check matrix can be partitioned into an array of block matrices, each block matrix is either a zero matrix or a right cyclic shift of an identity matrix. The parity check matrix designed in this way can be conveniently represented by a base (block) matrix. The main advantage is that they offer high throughput at low implementation complexity and they are used in many applications and communication standards. The ntLDPCD-Ghn core is fully compliant with various wireless and wireline communication standards including ITU-T G.9960 (G.hn), IEEE 802.16e (WiMAX), IEEE 802.11n/ac (WiFi) etc. The core implements an approximation of the log-domain LDPC iterative decoding algorithm, is highly reconfigurable and is able to support different sub-matrix sizes (Z) of LDPC-BC, that are tailored for specific applications. It also supports varying on the fly code rates, decoding iterations and input data width. The implementation is flexible, high speed, low cost and has a simple interface for easy integration in SoC applications.
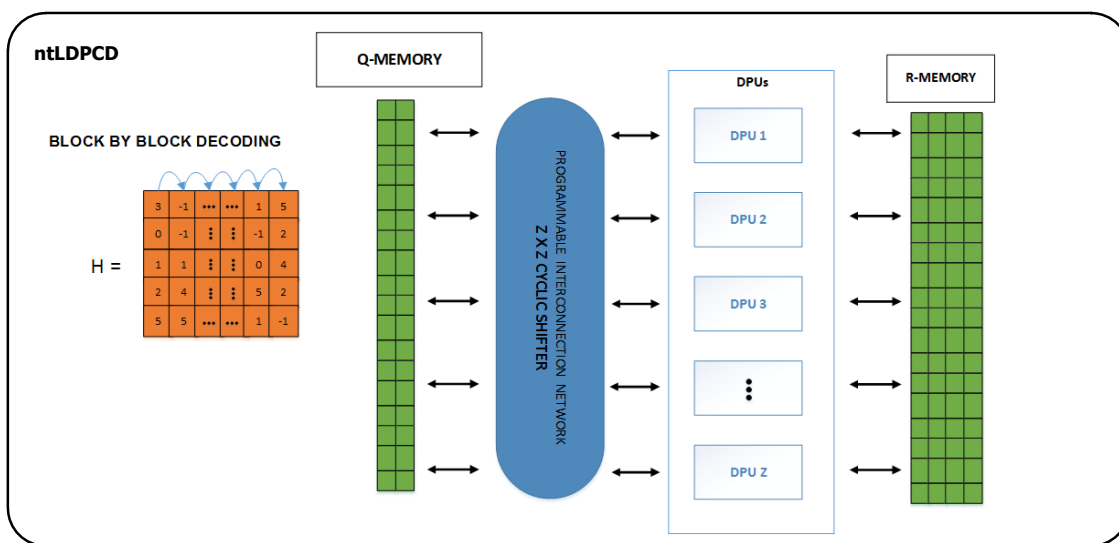
The BER vs SNR performance curves of the ntLDPCD-Ghn core for various iterations and for L-OMS algorithm compared to the SPA algorithm is shown below. As illustrated the implemented L-OMS algorithm presents almost the same coding gain with less iterations when compared to the SPA algorithm.

Forward Error Correction
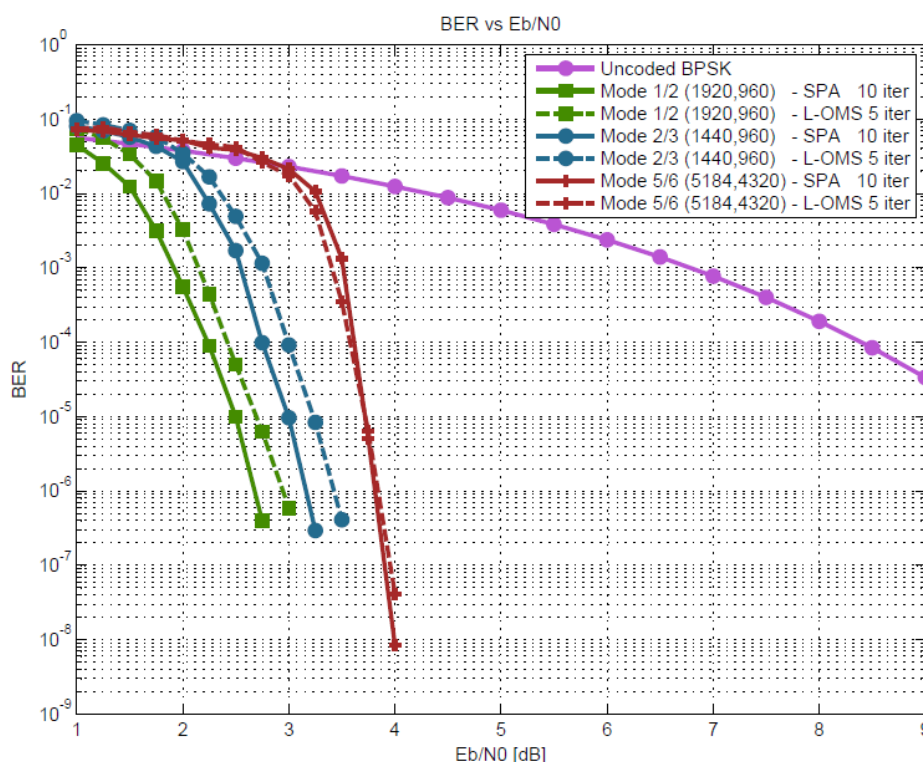
ntLDPCE-DVB-S2

DVB S2 Low Density Parity Check Encoder
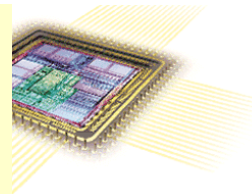
The ntLDPCE-DVB-S2 IP Core implements the encoding procedure for LDPC Block Codes (LDPC-BC) or QC-LDPC Quasi-Cyclic LDPC Codes compliant with the DVB-S2 standard. These LDPC codes are transformed to approximate block-structured LDPC codes with circular block matrices. The entire parity check matrix can be partitioned into an array of block matrices; each block matrix is either a zero matrix or a right cyclic shift of an identity matrix. The parity check matrix pre-processed in this way can be conveniently represented by a base matrix represented by cyclic shifts. The core is highly reconfigurable and it is able to support the sub-matrix size (Z=360) of QC-LDPC, that is tailored for the DVB-S2 standard.

The ntLDPCE-DVB-S2 Encoder has partial parallel architecture and supports Z=360 parallel input bits per clock cycle. The encoder receives information data, generates the parity bits and forms the codeword that will be transmitted. A selected mode requires K_LDPC/Z clock cycles in order to feed the encoder with a single block of information. The encoder architecture input is configured to support double buffering.
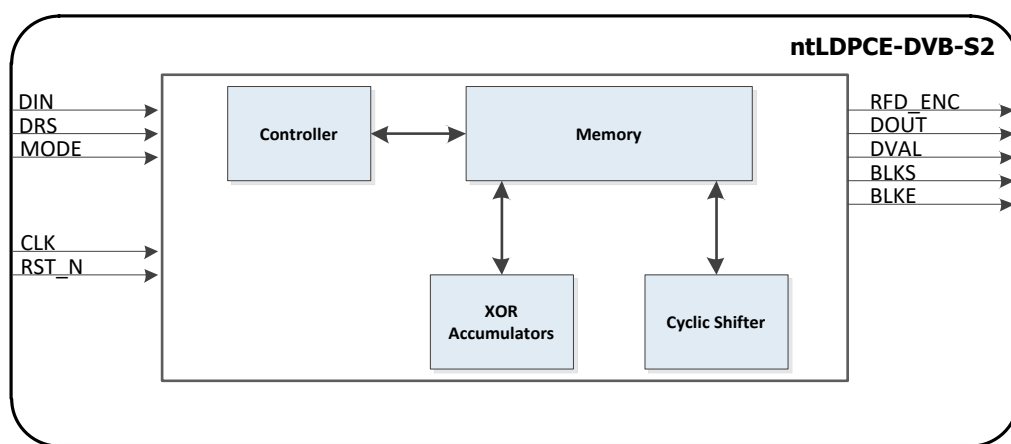
The following table demonstrates the DVB-S2 encoding modes of operation.

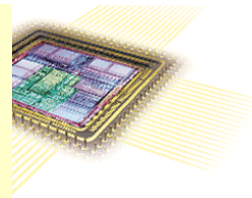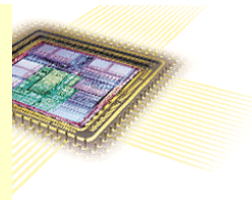| Mode	Codeword size	DVB S2 Mode (R=K/N)	Mode	Codeword size	DVB S2 Mode (R=K/N)
0	16200	DVB S2 1/4 (R=1/5)	11	64800	DVB S2 2/3
1	64800	DVB S2 1/4	12	16200	DVB S2 3/4 (R=11/15)
2	16200	DVB S2 1/3	13	64800	DVB S2 3/4
3	64800	DVB S2 1/3	14	16200	DVB S2 4/5 (R=7/9)
4	16200	DVB S2 2/5	15	64800	DVB S2 4/5
5	64800	DVB S2 2/5	16	16200	DVB S2 5/6 (R=2=37/45)
6	16200	DVB S2 1/4 (R=4/9)	17	64800	DVB S2 5/6
7	64800	DVB S2 1/2	18	16200	DVB S2 8/9
8	16200	DVB S2 3/5	19	64800	DVB S2 8/9
9	64800	DVB S2 3/5	20	-	RESERVED
10	16200	DVB S2 2/3	21	64800	DVB S2 9/10

The ntLDPCE-DVB-S2 core has been synthesized using Xilinx Vivado Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details and performance metrics of the ntLDPCE-DVB-S2 core configured for DVB S2 standard are shown in the tables below.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	DVB-S2 compliant	3615 FFs / 13485 LUTs / 222 Block RAMs	140

ntLDPCD-DVB-S2

DVB S2 Low Density Parity Check Decoder

The ntLDPCD-DVB-S2 IP Core implements an approximation of the log-domain LDPC iterative decoding algorithm (Belief propagation) known as Layered Offset Min-Sum Algorithm. As an alternative for better performance and error floor elimination for low code rates, the Layered Lambda-2 (λ=2) Min Algorithm has also been implemented, with a trade-off of increased hardware cost. Selection between the two algorithms is easily made via a generic value before synthesis. The core is highly reconfigurable and via a complex off-line preprocessing procedure, it is tailored for the DVB-S2 standard LDPC matrices.

The ntLDPCD-DVB-S2 IP Core has been implemented with partial block parallel architecture and can support a sub-matrix size Z=360 for structured-block LDPC codes. The decoder receives the distorted codeword from the demapper in LLRs and based on the maximum iterations specified produces the decoded result. The decoder supports input of Z=360 parallel Log-Likelihood Ratios (LLRs) per clock cycle. The representation of LLRs is notated as (wl,fr). The wl notation stands for wordlength bits including the sign bit and the fr notation accounts for the fractional bits. For this prototype implementation, the decoder is configured to receive LLRs with representation of (6,1) meaning that total word length is 6 bits, 1 bit is used for sign, 4 bits are used for the integer part and 1 bit is used for the fractional part. This feature is generic, algorithm dependent and can be adjusted based on the requirements of BER/FER performance, area cost and throughput requirements. Bit growth due to iterative decoding has also been considered and can be dynamically calibrated via another set of generics.

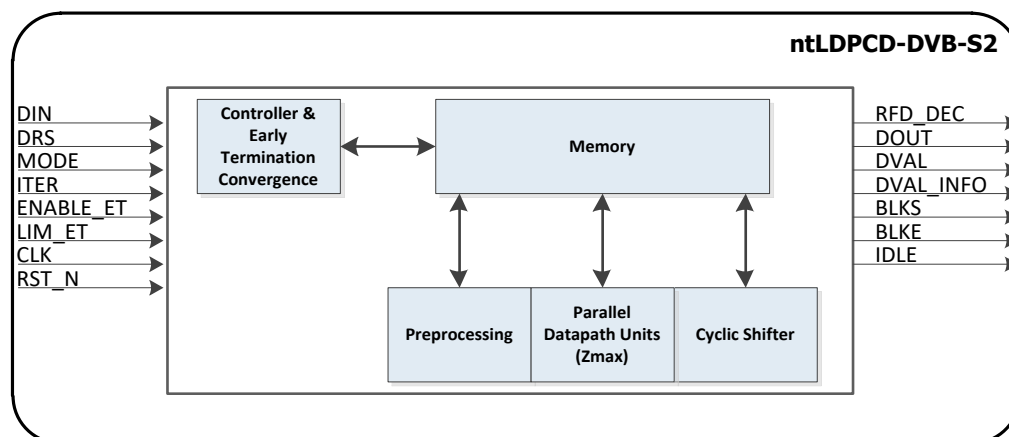An early termination mechanism has been installed and may be enabled by the 'enable_et' input port. When the mechanism is enabled and an internal criterion is met, the decoder controller will terminate the decoding process, before the maximum amount of iterations is performed and flush out the corrected codeword. When the termination criterion is met for lim_et times, then the early termination mechanism is activated.

Features

Features
Complex off-line LDPC matrices preprocessing for optimum RTL implementation efficiency.
Generic layered LDPC decoder architecture, that that can be tailored to implement any standard, thanks to Noesis Technologies patent pending off-line matrices preprocessing procedure.
Generic LLR input and internal fixed point precision.
Generic selection between Offset Min-Sum and Lambda 2 Min decoding Algorithms. Lambda 2 Min achieves even better performance and low code-rate error floors elimination at the expense of increased hardware utilization.
Early termination mechanism with robust convergence criterion for throughput increase without performance loss.
Competitive Frame Error Rate vs SNR Performance that meets the DVB-S2 standard's Quasi Error Free (QEF) requirements.

ntLDPCD-DVB-S2
DVB S2 Low Density Parity Check Decoder
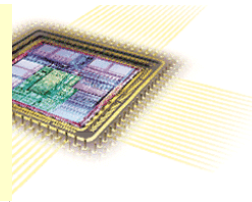
The ntLDPCD-DVB-S2 core has been synthesized using Xilinx Vivado Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details and performance metrics of the ntLDPCD-DVB-S2 core configured for DVB S2 standard are shown in the tables below.

Silicon Vendor	Device	Algorithm Config	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	Offset Min Sum	35204 FFs / 116805 LUTs / 320 Block RAMs	50
Xilinx	Kintex 7 XC7K410T-2	λ Min	37391 FFs / 131164 LUTs / 320 Block RAMs	50

The ntLDPCD-DVB-S2 achieves exceptional error correction performance as illustrated in the following FER vs Es/No graphs for all modes of operation as described in DVB S2 specification. The measurement conditions were as follows:

□ Input LLR S6.1 (signed, 4 integer bits, 1 fractional bit)
□ Approximate LLR calculation without proportional σ**2 scaling
□ AWGN channel impairments
□ QPSK modulation
□ Early Termination enabled after 4 converging iterations
□ Concatenation with outer 0.99 code rate block code

ntLDPCE-DVB-S2X
DVB S2X Low Density Parity Check Encoder
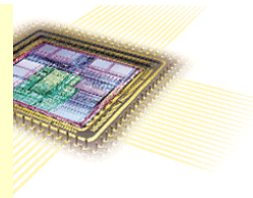
The ntLDPCE-DVB-S2X IP Core implements the encoding procedure for LDPC Block Codes (LDPC-BC) or QC-LDPC Quasi-Cyclic LDPC Codes for the DVB S2X standard. These LDPC codes are transformed to approximate block-structured LDPC codes with circular block matrices. The entire parity check matrix can be partitioned into an array of block matrices; each block matrix is either a zero matrix or a right cyclic shift of an identity matrix. The parity check matrix pre-processed in this way can be conveniently represented by a base matrix represented by cyclic shifts. The core is highly reconfigurable and it is able to support the sub-matrix size (Z=360) of QC-LDPC, that is tailored for the DVB S2X standard.

The ntLDPCE-DVB-S2X encoder has partial parallel architecture and supports Z=360 parallel input bits per clock cycle. The encoder receives information data, generates the parity bits and forms the codeword that will be transmitted. A selected mode requires K_LDPC/Z clock cycles in order to feed the encoder with a single block of information. The encoder architecture input is configured to support double buffering.

The ntLDPCE-DVB-S2X encoder may optionally include all DVB-S2 modes of operation. Optional DVB S2X puncturing and/or shortening is implemented according to the following specification, wherever it applies according to the DVB S2X standard. ntLDPCE-DVB-S2X encoder expects (kldpc+Xs) bits at its input in 360-bit code-words and forces the first Xs input bits to zero internally before encoding. After encoding the IP outputs the whole (16200,32400,64800) encoded block in 360-bit code-words. Additionally it calculates puncturing pattern internally and indicates via the 360-bit output mask (MASKO) which exactly bits need to be omitted (either Xs or Xp). Xs, kldpc, P and Xp are defined identically to the DVB S2X standard.
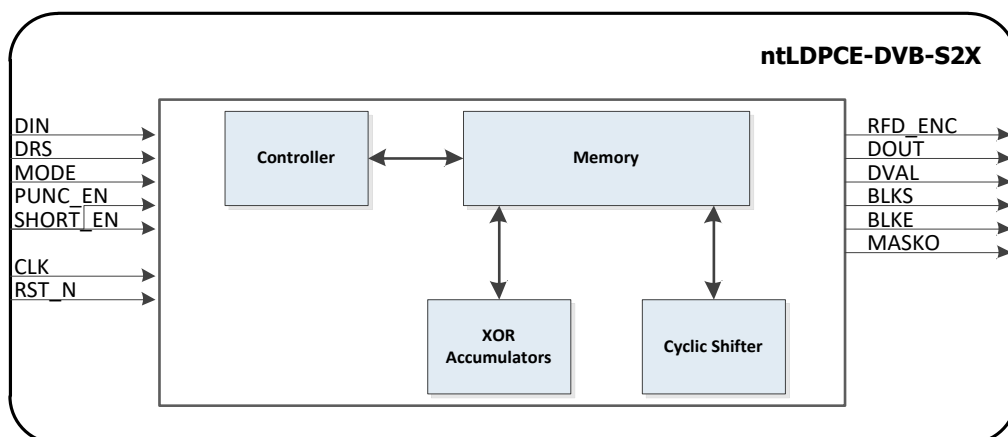
The following table demonstrates the DVB-S2X encoding modes of operation.

Mode	Codeword size	DVB S2X Mode (R=K/N)	Mode	Codeword size	DVB S2X Mode (R=K/N)
0	16200	DVB S2X 11/45	18	64800	DVB S2X 26/45
1	16200	DVB S2X 4/15	19	64800	DVB S2X 18/30
2	16200	DVB S2X 14/45	20	64800	DVB S2X 28/45
3	16200	DVB S2X 7/15	21	64800	DVB S2X 23/36
4	16200	DVB S2X 8/15	22	64800	DVB S2X 116/180
5	16200	DVB S2X 26/45	23	64800	DVB S2X 20/30
6	16200	DVB S2X 32/45	24	64800	DVB S2X 124/180
7	32400	DVB S2X 1/5	25	64800	DVB S2X 25/36
8	32400	DVB S2X 11/45	26	64800	DVB S2X 128/180
9	32400	DVB S2X 1/3	27	64800	DVB S2X 13/18
10	64800	DVB S2X 2/9	28	64800	DVB S2X 132/180
11	64800	DVB S2X 13/45	29	64800	DVB S2X 22/30
12	64800	DVB S2X 9/20	30	64800	DVB S2X 135/180
13	64800	DVB S2X 90/180	31	64800	DVB S2X 140/180
14	64800	DVB S2X 96/180	32	64800	DVB S2X 7/9
15	64800	DVB S2X 11/20	33	64800	DVB S2X 154/180
15	64800	DVB S2X 100/180	34-55	-	Optional DVB-S2 modes support
17	64800	DVB S2X 104/180			

ntLDPCE-DVB-S2X

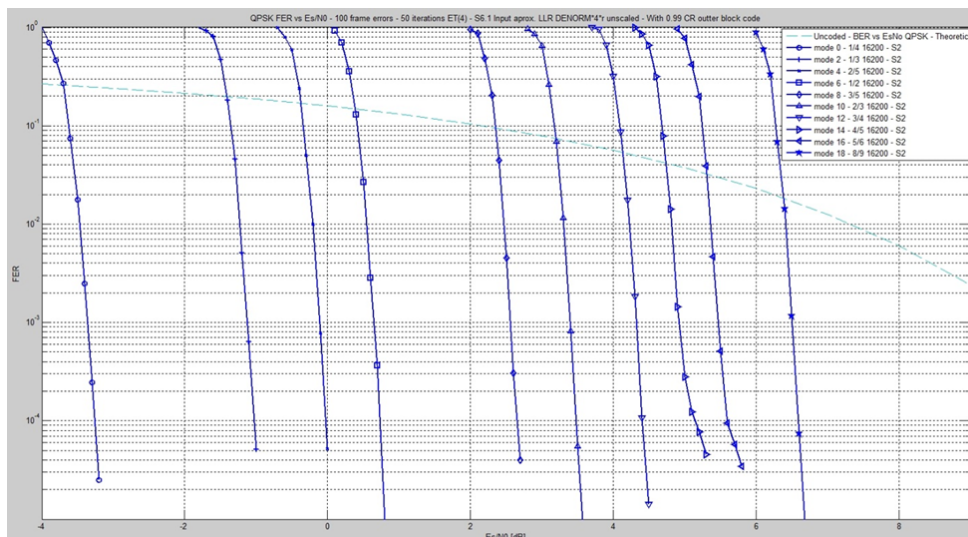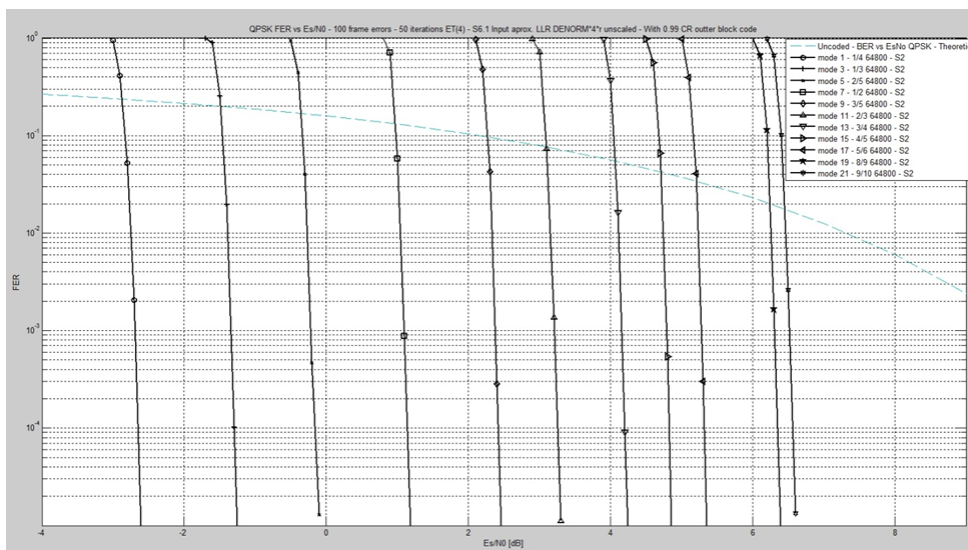## DVB S2X Low Density Parity Check Encoder
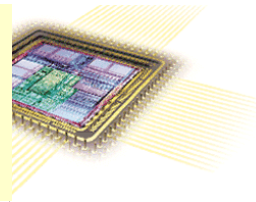
The ntLDPCE-DVB-S2X core has been synthesized using Xilinx Vivado Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details and performance metrics of the ntLDPCE-DVB-S2X core configured for DVB S2X standard are shown in the tables below.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	DVB-S2X compliant	3615 FFs / 13485 LUTs / 222 Block RAMs	140

ntLDPCD-DVB-S2X

DVB S2X Low Density Parity Check Decoder

The ntLDPCD-DVB-S2X IP Core implements an approximation of the log-domain LDPC iterative decoding algorithm (Belief propagation) known as Layered Offset Min-Sum Algorithm. As an alternative for better performance and error floor elimination for low code rates, the Layered Lambda-2 (λ=2) Min Algorithm has also been implemented, with a trade-off of increased hardware cost. Selection between the two algorithms is easily made via a generic value before synthesis. The core is highly reconfigurable and via a complex off-line preprocessing procedure, it is tailored for the DVB-S2X standard LDPC matrices.

The ntLDPCD-DVB-S2X IP Core has been implemented with partial block parallel architecture and can support a sub-matrix size Z=360 for structured-block LDPC codes. The decoder receives the distorted codeword from the demapper in LLRs and based on the maximum iterations specified produces the decoded result. The decoder supports input of Z=360 parallel Log-Likelihood Ratios (LLRs) per clock cycle. The representation of LLRs is notated as (wl,fr). The wl notation stands for wordlength bits including the sign bit and the fr notation accounts for the fractional bits. For this prototype implementation, the decoder is configured to receive LLRs with representation of (6,1) meaning that total word length is 6 bits, 1 bit is used for sign, 4 bits are used for the integer part and 1 bit is used for the fractional part. This feature is generic, algorithm dependent and can be adjusted based on the requirements of BER/FER performance, area cost and throughput requirements. Bit growth due to iterative decoding has also been considered and can be dynamically calibrated via another set of generics.

An early termination mechanism has been installed and may be enabled by the 'enable_et' input port. When the mechanism is enabled and an internal criterion is met, the decoder controller will terminate the decoding process, before the maximum amount of iterations is performed and flush out the corrected codeword. When the termination criterion is met for lim_et times, then the early termination mechanism is activated. The ntLDPCD-DVB-S2X decoder may optionally include all DVB-S2 modes of operation.

Features

Complex off-line LDPC matrices preprocessing for optimum RTL implementation efficiency.
Generic layered LDPC decoder architecture, that that can be tailored to implement any standard, thanks to Noesis Technologies patent pending off-line matrices preprocessing procedure.
Generic LLR input and internal fixed point precision.
Generic selection between Offset Min-Sum and Lambda 2 Min decoding Algorithms. Lambda 2 Min achieves even better performance and low code-rate error floors elimination at the expense of increased hardware utilization.
Early termination mechanism with robust convergence criterion for throughput increase without performance loss.
Competitive Frame Error Rate vs SNR Performance that meets the DVB-S2 standard's Quasi Error Free (QEF) requirements.
IP may also be backwards compatible with all DVB-S2 LDPC modes.

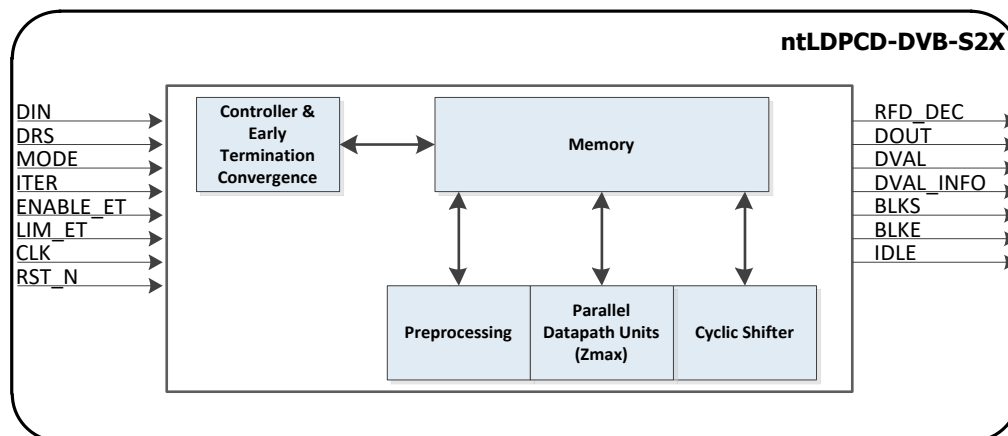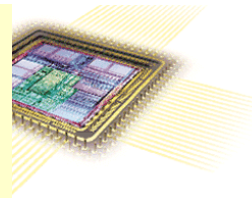# ntLDPCD-DVB-S2X
DVB S2X Low Density Parity Check Decoder
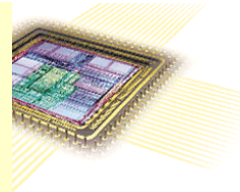
The ntLDPCD-DVB-S2X core has been synthesized using Xilinx Vivado Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details and performance metrics of the ntLDPCD-DVB-S2X core configured for DVB S2X standard are shown in the tables below.

Silicon Vendor	Device	Algorithm Config	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	Offset Min Sum	35204 FFs / 116805 LUTs / 320 Block RAMs	50
Xilinx	Kintex 7 XC7K410T-2	λ Min	37391 FFs / 131164 LUTs / 320 Block RAMs	50

The ntLDPCD-DVB-S2X achieves exceptional error correction performance as illustrated in the following FER vs Es/No graphs for all modes of operation as described in DVB S2X specification. The measurement conditions were as follows:

  □ Input LLR S6.1 (signed, 4 integer bits, 1 fractional bit)
  □ Approximate LLR calculation without proportional σ**2 scaling
  □ AWGN channel impairments
  □ QPSK modulation
  □ Early Termination enabled after 4 converging iterations
  □ Concatenation with outer 0.99 code rate block code
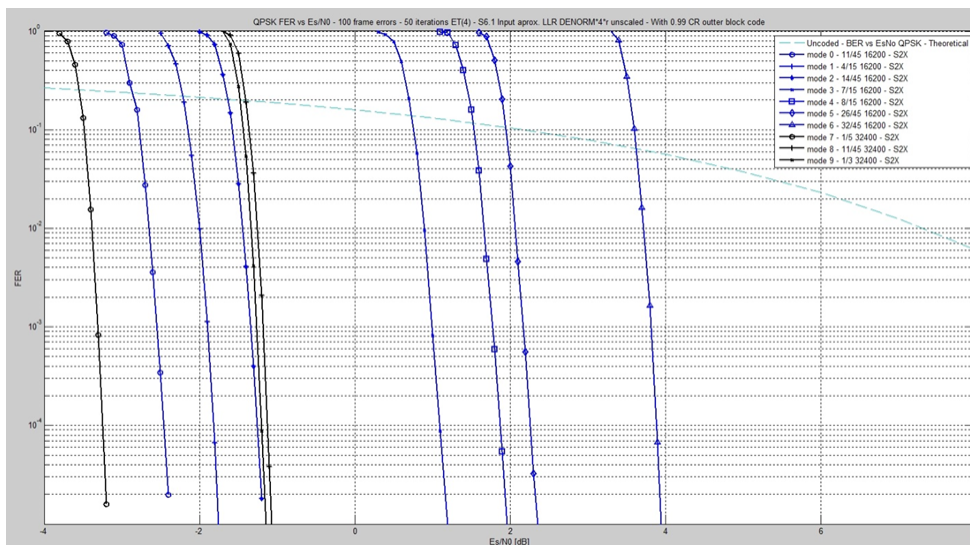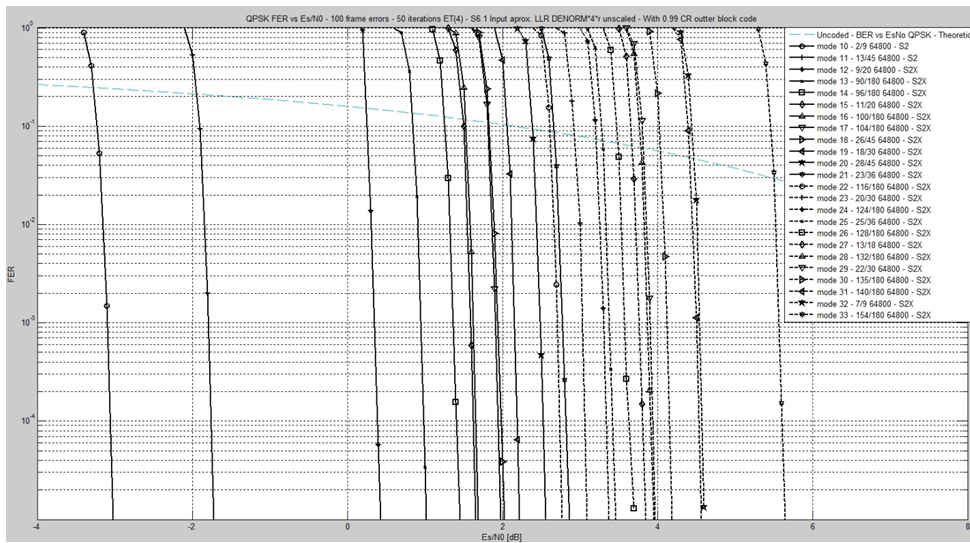
ntINT_DEINT
Fully Configurable Interleaver—Deinterleaver

Error detection and correction are perhaps the most important quality factors to observe when evaluating a digital transmission system. A system's noise environment can cause errors in the transmitted message degrading the credibility of the system. Digital communications systems designers can choose among many types of error-correction codes to reduce the effect of errors in stored or transmitted data. Most common error correcting codes are designed to correct random errors i.e. errors that are independent of each other and distributed uniformly in time. However errors that occur in bursts ,i.e. errors that occur sequentially in time and as groups, tend to be problematic for most FEC schemes. Block codes, and in particular Reed-Solomon codes, can handle burst errors effectively only as long as the number of errors per data block is below a certain limit. Interleaving is a simple, yet powerful technique that can be used to extend the error correcting capability of a Reed-Solomon code and other FEC codes. The ntINT_DEINT interleaver/de-interleaver subsystem rearranges the encoded symbols over multiple data blocks. This effectively spreads out long burst noise sequences so they appear to the decoder as independent random symbol errors or shorter more manageable burst errors. This is achieved by using the interleaving function that changes the order of data before the transmission on the channel so that any adjacent symbols are well separated during transmission. The symbols are then reordered by the deinterleaving function during reception. Block and Convolutional are the most frequently used interleaver types.

Features

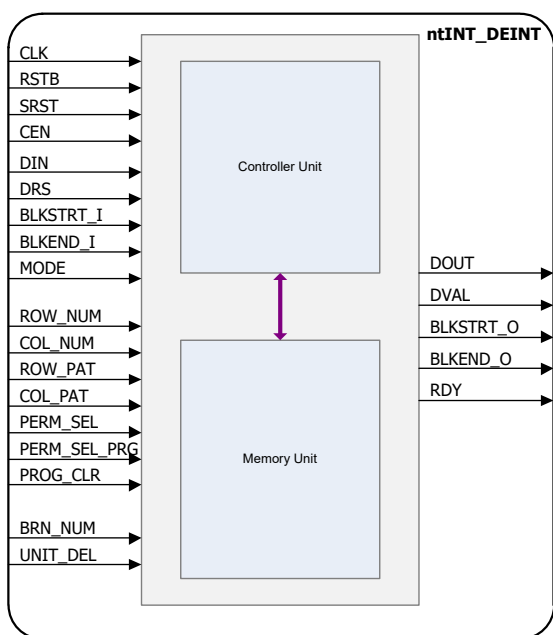| Features |
|---|
| Fully configurable, convolutional and rectangular interleaver / deinterleaver. |
| Compliant to a variety of industry standards such as DVB, ATSC, IEEE 802.16. etc. |
| Rectangular Block (de) interleaver configuration:
- Block size
- Number of rows
- Number of columns
- Rows and/or columns permutations |
| Convolutional (de) interleaver configuration:
- Number of branches
- Configurable branch length |
| Supports continuous block data flow. |
| Configurable number of bits per symbol. |
| Handshaking logic for I/O data flow control. |
| Fully synchronous design, using single clock. |
| Silicon proven in ASIC and Xilinx FPGA implementation technologies for a variety of applications. |

The ntINT_DEINT core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	Convolutional / NUMBITS=8 MAX_BRN_NUM=4 / MAX_UNIT_DEL=4	253 Slices / 3 Block RAMs	157
Xilinx	Kintex 7 XC7K410T-2	Convolutional / NUMBITS=16 MAX_BRN_NUM=16 / MAX_UNIT_DEL=16	744 Slices / 13 Block RAMs 1 DSP48E1	132
Xilinx	Kintex 7 XC7K410T-2	Convolutional / NUMBITS=8 MAX_BRN_NUM=32 / MAX_UNIT_DEL=32	1276 Slices / 31 Block RAMs 1 DSP48E1	124
Xilinx	Kintex 7 XC7K410T-2	Block / NUMBITS=8 / MAX_ROW_NUM=32 MAX_COL_NUM=32 / MAX_PREM_NUM=4	185 Slices / 4 Block RAMs	261
Xilinx	Kintex 7 XC7K410T-2	Block / NUMBITS=16 / MAX_ROW_NUM=64 MAX_COL_NUM=64 / MAX_PREM_NUM=4	246 Slices / 6 Block RAMs	226
Xilinx	Kintex 7 XC7K410T-2	Block / NUMBITS=64 / MAX_ROW_NUM=128 MAX_COL_NUM=128 / MAX_PREM_NUM=4	320 Slices / 66 Block RAMs 1 DSP48E1	166

Voice & Data Compression
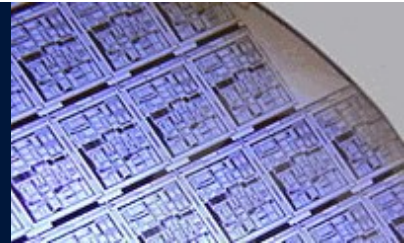
Voice compression technology is widely used in digital communication systems such as wireless systems, VoIP, and video conference technology. Voice compression reduces data redundancy and thus eases bandwidth requirements. The International Telecommunication Union (ITU) has standardized a number of speech compression algorithms for a variety of compression rates as well as Mean Opinion Scores (MOS). Noesis Technologies provides a series of silicon IPs of the most popular voice codecs (G711, G726, G729, CVSD), providing compression rates ranging from 64 kbps down to 8 kbps.

In addition, Noesis Technologies offers a proprietary implementation of Huffman block differential lossless data compression algorithm. This core is ideal for use in low power applications like Wireless Sensor Networks (WSN) as well as any other application with slow changing nature of data, to fully benefit from the differential nature of the algorithm.

ntG711
A/u Law Codec— ITU-T G711 compliant

The ntG711 core implements the ITU G.711 compliant compressing and expanding functions. It is comprised of a compressor and an expander unit. The compressor unit performs compression of the 16-bit uniform PCM to a 8-bit A/μ-law word. The expander unit decompresses the 8-bit A/μ-law to 14-bit uniform PCM word. The ntG711 core is programmable and its functionality is controlled by the following control bits.

law : This bit selects the coding rule to be used. When '0' μ-law is selected, when '1' A-law is selected.
A_inv_dis : This bit activates/disactivates the inversion of even bits of the input word for the A-law case.
u_inv_dis : This bit activates/disactivates the inversion of bits of the input word for the μ-law case.
comp_dis : This bit selects the representation format of the output vector. When '0' is in 2's complement format, when '1' is in sign magnitude format.
The ntG711 core can be used in a variety of applications, including PCM codecs, voice compression and expanding as well as a front–end for any DSP processing of 64 kbps voice.

The PCM to A-law and to μ-law transformation

A/μ-LawCompressor
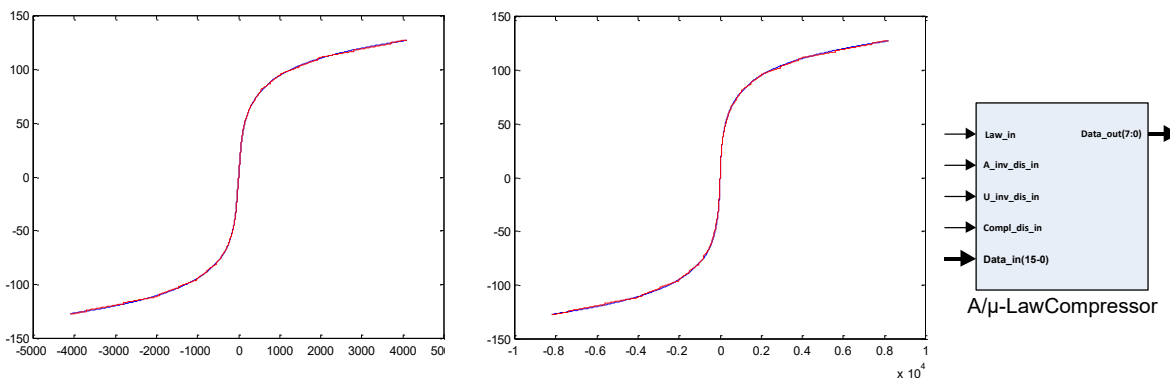
A-law and μ-law to PCM transformation

A/μ-LawExpander
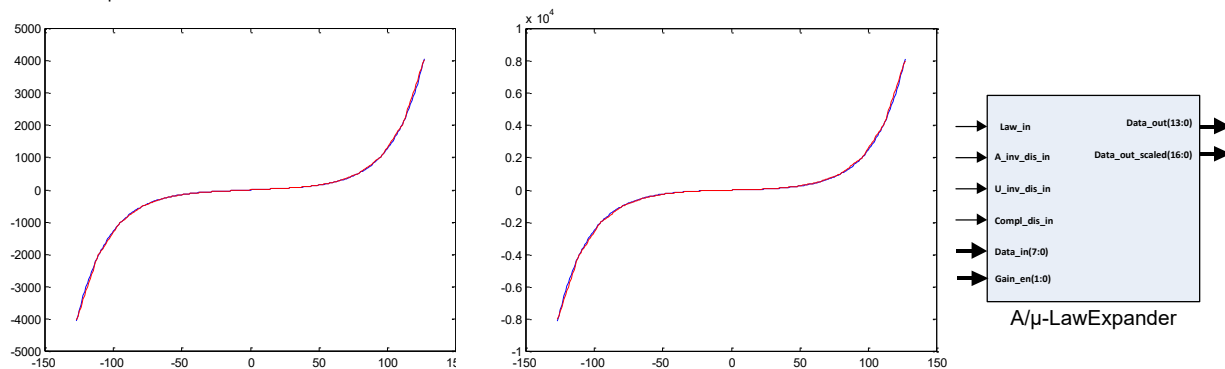
The ntG711 core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Virtex-5	54 CLB Slices (Compressor unit) 75 CLB Slices / 1 DSP48E (Expander unit)	115 MHz
TSMC	180 nm	610 gates (Compressor unit) 720 gates (Expander Unit)	

ntG726

Multi Channel ADPCM Codec — ITU-T G726 compliant
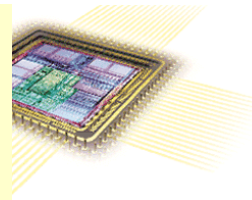
The ntG726 core is fully compliant with G.726 standard and supports up to 64 full duplex voice channels. The G.726 recommendations specifies the conversion of a 64 kbps A-law or μ-law pulse code modulation (PCM) to and from a 40, 32, 24 and 16 kbps channel. This conversion is applied to the PCM bit stream using an ADPCM transcoding technique. The ntG726 core can be configured 'on-the-fly' for A-law or μ-law linear code and conversion rate on a per channel basis. The core is used in applications that require reduction in transport and storage bandwidth requirements. It significantly offloads CPU tasks as a co-processing system element.
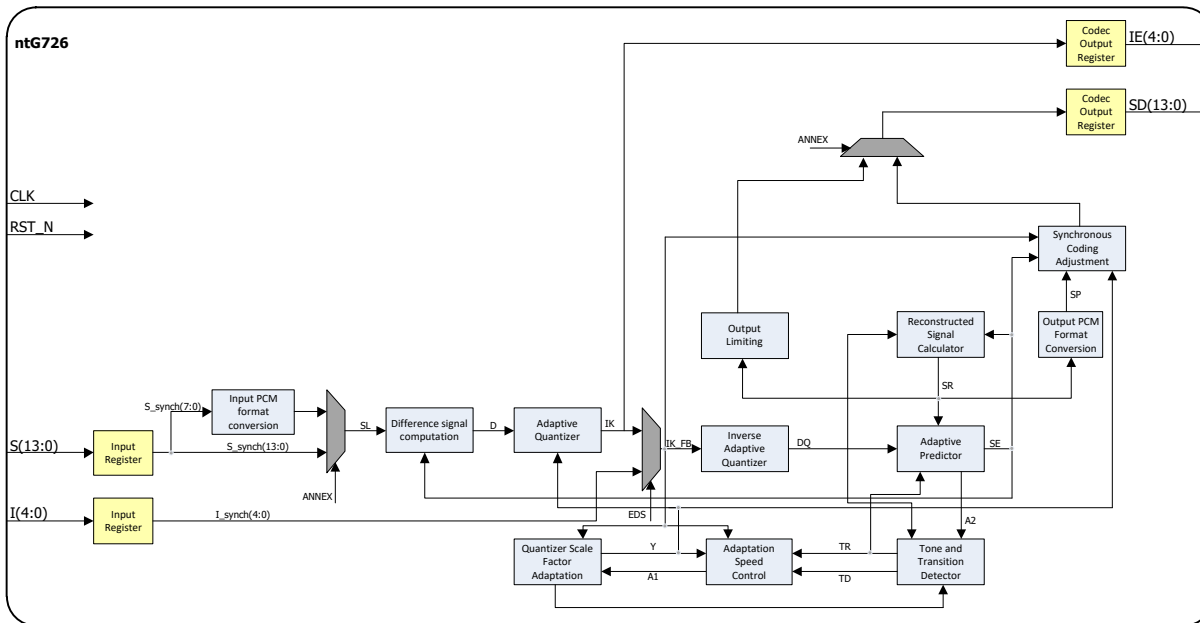
Features

Features
Compliant with ITU G.721, G.723, G.726 and G.726-Annex recommendations.
'On-the-fly' configuration for variable compression rate, PCM law.
Process capability of up to 64 full duplex or up to 128 half duplex voice channels.
Burst and continuous mode support.
No register based configuration is required.
A-law, μ-law linear code format selection.
Fully comforted to ITU test vectors (ITUG.726-A2).
Fully synchronous design, using single clock.
Portable to any FPGA/ASIC technology.

The ntG726 core has been targeted to both ASIC and FPGA technologies for various applications. Noesis Technologies can also deliver netlist versions of the core optimized to specific area resources and performance requirements.
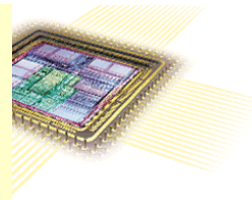
Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Virtex-II	2515 CLB Slices	60
TSMC	180 nm	24 K gates	200

ntG729

Multi Channel CS-ACELP 8 kbps Codec — ITU-T G729A compliant

The growth in wireless communication systems, cellular mobile radio and VoIP technology has created the imperative need for bandwidth efficient, high speed quality voice coding algorithms. The ITU-T G.729 CS-ACELP is a high speech quality, low-bit rate (8kbps) codec that has been proposed to meet the voice compression requirements of a modern communication system. However the real time SW implementation of a multi-channel ITU-T G.729 compliant voice codec in conventional DSP processors is prohibitive due to the intensive amount of signal processing power required by the algorithm. To overcome this limitation Noesis Technologies has developed a revolutionized, highly efficient hybrid architecture that implements real time multi-channel G729A voice coding and exhibits the best performance-silicon area ratio available in the industry. The ntG729 IP Core can be used as a coprocessor to any processor type and can save significant computing resources for the main processor by efficiently executing the computationally intensive speech coding G729A algorithmic operations.
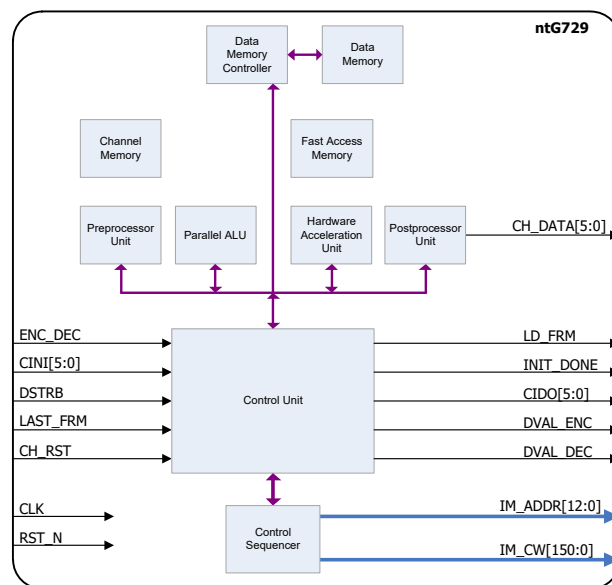
Features

Features
Voice codec capable of multi-channel 8kbps voice compression based on ITU-T G729A standard.
Selective Channel initialization.
AMBA bus support for easy SoC integration.
Outperforms by a factor of two by supporting twice as many voice channels.
It requires only half the gate count thus increasing power efficiency and reducing costs.
Best performance/silicon area ratio available in the industry.
Fully synchronous design, using single clock.
Portable to any FPGA/ASIC technology.
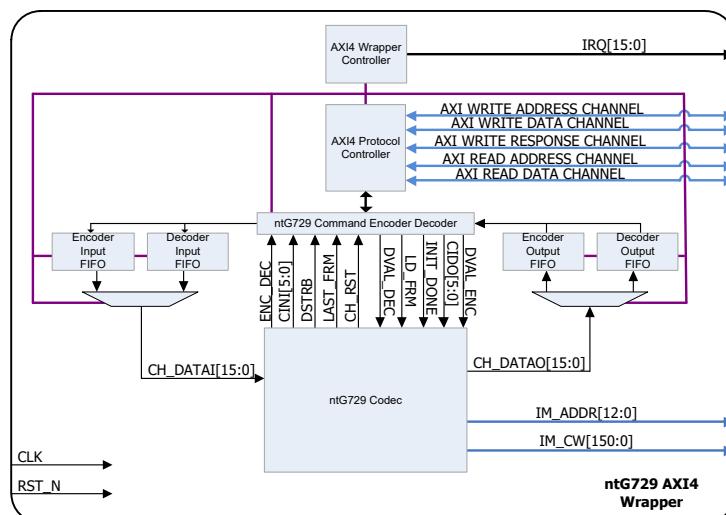
The ntG729 core has been core has been implemented to Xilinx devices.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Virtex-7	18K CLB Slices / 31 Block RAMs / 105 DSP Slices	117

ntCVSD

Continuously variable slope delta modulation

The fundamental principle of the CVSD algorithm is the encoding of one bit per sample. For example an audio signal sampled at 32 KHz will be compressed to 32 Kbps.

The ntCVSD codec IP core can be configured to operate either as an encoder or as a decoder functional block. In encoder mode the core accepts input data at a rate of 8 KHz/128 Kbps or 16KHz/256Kbps and are sampled when the data strobe signal is asserted high. Higher input sampling data rates can also be supported with no up-sampling provision. The sampled input data can either be initially up-sampled to 64 KHz by using an interpolation filter in order to improve speech/audio quality before entering into the actual CVSD codec unit or can be just directly fed into the CVSD codec unit with no previous processing. These samples are then driven to a digital comparator in order to be compared with a reference signal value. If the input sample is greater than the reference signal then a logic 1 is transmitted and a step value is added on the reference signal. If the input sample is less than the reference signal then a logic 0 is transmitted and a step value is subtracted from the reference signal. The transmitted bits are also stored in an N-bit shift register. Depending on the shift register contents, a decision is made whether a slope overload has occurred and the step value is adjusted accordingly in order to keep up with the changing slope of the input waveform. Depending on both the digital comparator and the slope overload decisions, an integrator estimation is generated in order to approximate the previous input value and drive it back to the digital comparator.

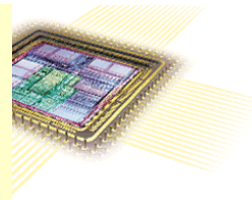In decoder mode the core accepts the compressed bit-stream and the incoming bits are sampled when the data strobe signal is asserted high. The received bit-stream is fed directly to the N-bit shift register and depending on the shift register contents, a decision is made whether a slope overload has occurred. The uncompressed signal is reconstructed through the integrator unit. The reconstructed output will be either be fed into the decimation filter to be down-sampled at the original sampling rate and then driven at the output or will be just directly driven at the output with no previous processing.

The ntCVSD codec is designed to compress 16-bit PCM speech/audio data for transmission in telecom networks or to decompress a received CVSD encoded bit-stream. It is compatible with legacy CVSD implementations and fully compliant with the Bluetooth CVSD specification.
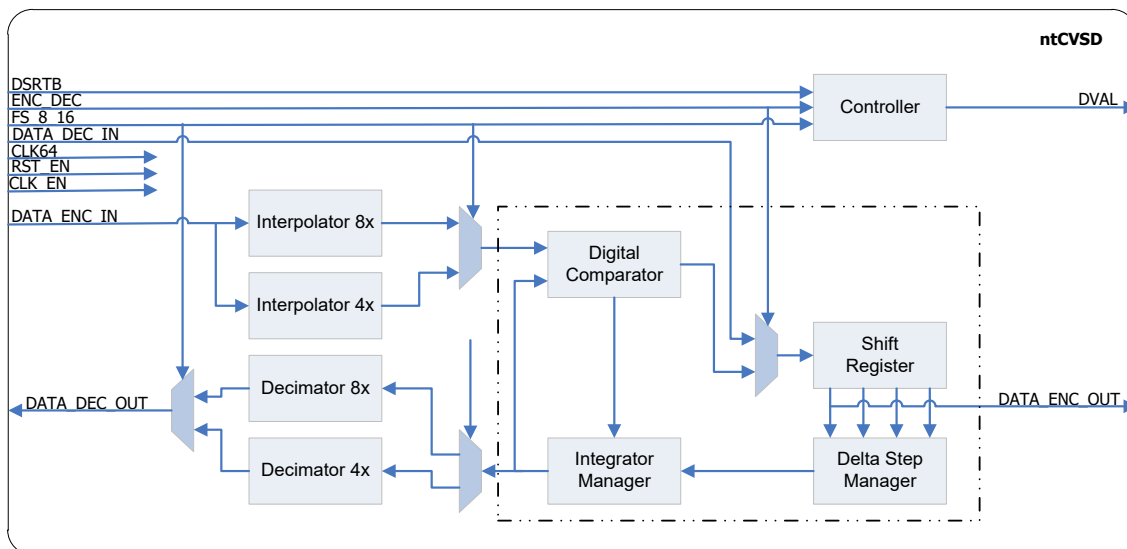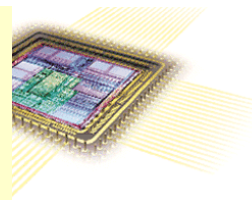

The ntCVSD core has been core has been implemented to Xilinx devices.

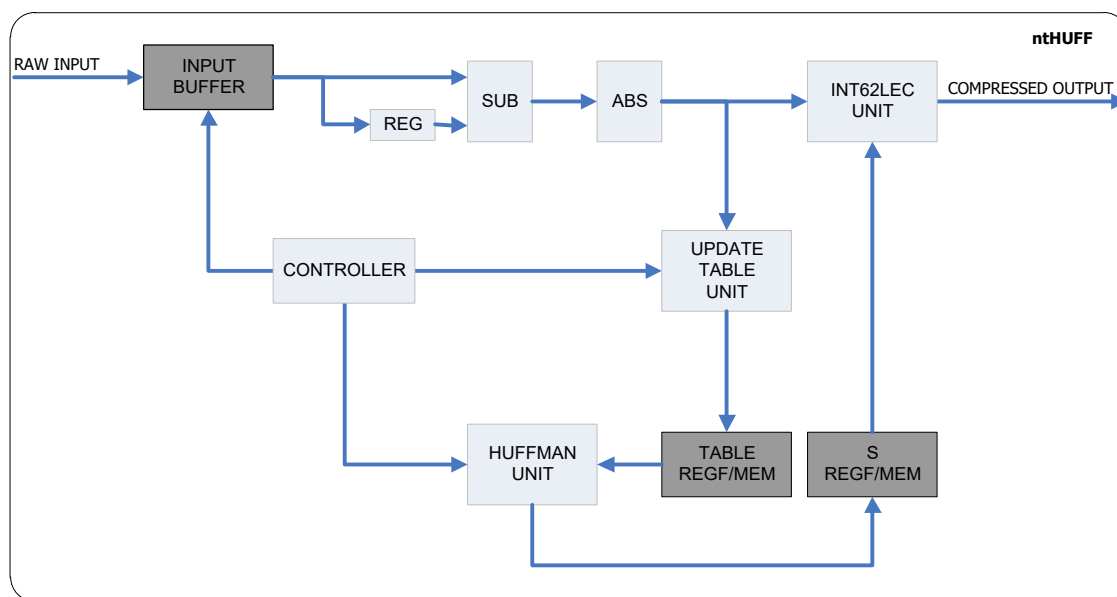| Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Virtex-6	349 CLB Slices / 31 DSP Slices	95

ntHUFF
Huffman algorithm compression engine

The ntHUFF compression module implements the Huffman Block differential compression algorithm. The core processes data blocks of 500 16-bit input samples "on the fly" with latency as little as 4 clock cycles. A small input buffer of configurable size stores incoming 16-bit samples and propagates them to the compression module when instructed by the local controller. Samples are propagated through the differential data path comprised of a subtractor and an absolute calculation unit. The absolute value of all samples is used to update a metric table with statistical information and is also used to produce the compressed output. This is the initialization phase of the system.

When samples equal to the defined block size have been collected, the controller enters calculations phase and pauses further samples propagation to the rest of the system. The Huffman microprocessor unit calculates and produces the Huffman (S) table, based on the populated metric table, which will be applied on the next block of incoming data. The custom microprocessor functions with encoded operations designed to optimize this phase. The core of the Huffman algorithm is implemented by performing parallel memory accesses on a parallel memory. A 512x36 instruction memory drives the micro-processor to execute all real time Huffman algorithm calculations. Worst case processing latency due to iterative algorithm nature is calculated to 4175 clock cycles per block of 500 samples.

Once the Huffman (S) table has been calculated, the controller resumes samples propagation through the differential data path and S is applied to the next block of incoming samples in order to produce the compressed output. When samples equal to the defined block size have been collected, the calculations phase is activated again and so on. A flush enable input port is provided, which applies zero padding on the last 32bit compressed output at the end of each data block.

The ntHUFF IP Core can be used in wireless sensor networks, medical applications as well as any data compression application with slow changing nature of data, to fully benefit from the differential nature of the algorithm.

The ntHUFF core has been core has been implemented to Xilinx devices.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3	1724 CLB Slices / 8 Block RAMs	63

Security
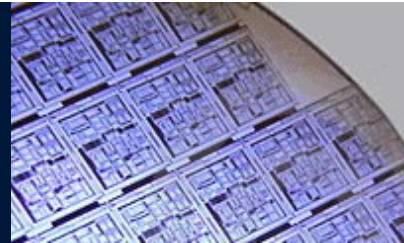
Noesis Technologies offers a range of security solutions ensuring privacy and authentication in digital transmissions. The Advanced Encryption Standard (AES) has been ratified in November 2001 by the National Institute of Standards and Technology (NIST) as the new encryption standard (FIPS PUB 197) to replace the existing, aging and vulnerable Data Encryption Standard (DES). An algorithm called Rinjdael has been finally selected for AES among a number of other candidates, after successfully meeting a set of criteria including not only security but performance and implementation feasibility in a variety of applications. Noesis Technologies has been involved in the development of cryptographic solutions for telecom and defense sector since 2000 and its class leading solutions have been silicon proven in multiple applications. A wide range of AES cores have been designed to support different performance & silicon area combinations that provide the optimum implementation for end-user application requirements. Due to a unique Galois Field Multiplier architectural implementation, the structural datapath element of all AES cryptographic engines, and the highly efficient algorithmic mapping techniques the family of ntAES IP cores exhibit the best performance-silicon area ratio available in the industry. In addition, Noesis Technologies provides SHA 256-bit authentication and RC4 encryption IP cores for a various speed & silicon area application requirements.

ntAES8
AES Low Power Encryption/Decryption Engine

ntAES8 core implements the NIST FIPS-197 Advanced Encryption Standard and can be programmed to either encrypt or decrypt 128-bit blocks of data using a 128-bit, 192-bit or 256-bit key. The ntAES8 has been carefully designed to require minimum logic resources rendering it an ideal solution for low power applications. This has been achieved by using an 8-bit data path size which means that 16 clock cycles are required to load/unload the 128-bit plaintext/ciphertext block. The encryptor receives the 128-bit plaintext block in 8-bit input symbols and generates the corresponding 128-bit ciphertext block in 8-bit output symbols using a supplied 128, 192, or 256-bit AES key. The pre-computed key values are read from an internal round key RAM. A key expander module is provided as an optional module to allow automatic generation and loading of the round key RAM. The decryptor implements the reverse function, generating plaintext from supplied ciphertext, using the same AES key as was used for encryption. The implementation is very low on latency, high speed with a simple interface for easy integration in SoC applications. The ntAES8 core can be used in a variety of applications, including:
- Electronic financial transactions.
        eCommerce, Banking, Securities exchange, Point-of-Sale
- Secure communications.
        Storage Area Networks (SAN), Virtual Private Networks (VPN) Video Conferencing, Voice services
- Secure environments.
        Satellite communications, Surveillance systems, Network appliances
- Personal mobile communications.
        Video phones, PDA, Point-to-Point Wireless

Key size	Throughput rate (Xilinx Spartan-3)
128 bits	53.3 Mbps
192 bits	44 Mbps
256 bits	37.4 Mbps

The ntAES8 core has been core has been implemented to Xilinx devices as well as TSMC ASIC libraries.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3	160 CLB Slices / 1 Block RAM	200
TSMC	180 nm	1226 gates / 7680 RAM bits	515

ntAES32
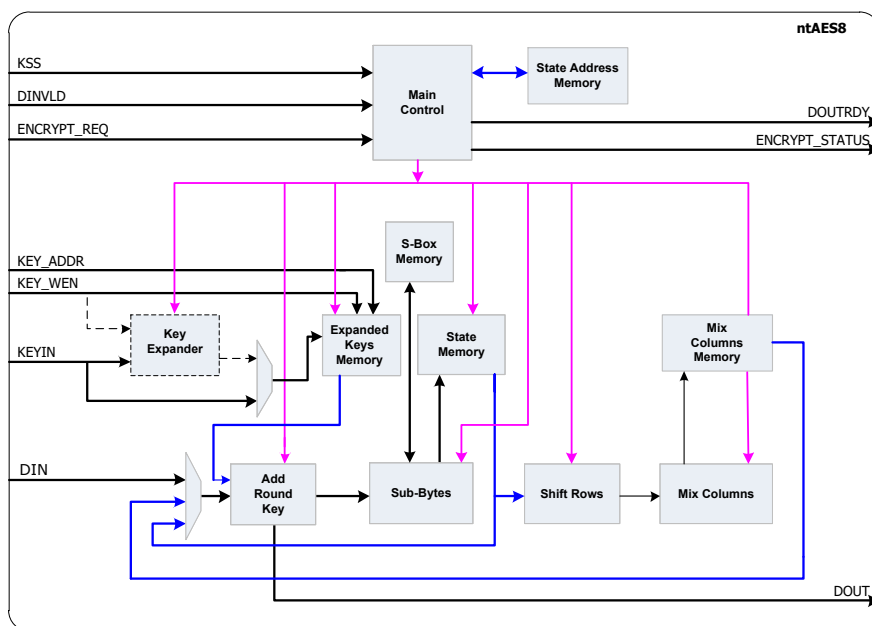## AES High Speed Encryption/Decryption Engine
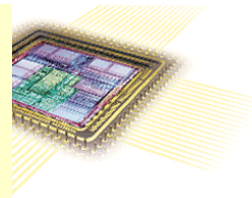
ntAES32 core implements the NIST FIPS-197 Advanced Encryption Standard and can be programmed to either encrypt or decrypt 128-bit blocks of data using a 128-bit, 192-bit or 256-bit key. The ntAES32 has been carefully designed for high throughput applications with optimal logic resources utilization. The encryptor core accepts a 128-bit plaintext input word, and generates a corresponding 128-bit ciphertext output word using a supplied 128, 192, or 256-bit AES key. The decryptor core provides the reverse function, generating plaintext from supplied ciphertext, using the same AES key as was used for encryption. The hardware roundkey expansion logic has been designed as a discrete building block. This allows either to build a complete stand-alone AES solution, or to save logic resources by leaving the key generation process to the user. Alternatively, the roundkey expansion logic can be shared between multiple encryption/decryption cores for optimal silicon area resources utilization. The implementation is very low on latency, high speed with a simple interface for easy integration in SoC applications.

The ntAES32 core can be used in a variety of applications, including:
- Electronic financial transactions.
    eCommerce, Banking, Securities exchange, Point-of-Sale
- Secure communications.
    Storage Area Networks (SAN), Virtual Private Networks (VPN) Video Conferencing, Voice services
- Secure environments.
    Satellite communications, Surveillance systems, Network appliances
- Personal mobile communications.
   Video phones, PDA, Point-to-Point Wireless

Key size	Throughput rate (Xilinx Virtex-5)
128 bits	550 Mbps
192 bits	460 Mbps
256 bits	400 Mbps

The ntAES32 core has been core has been implemented to Xilinx FPGA devices.

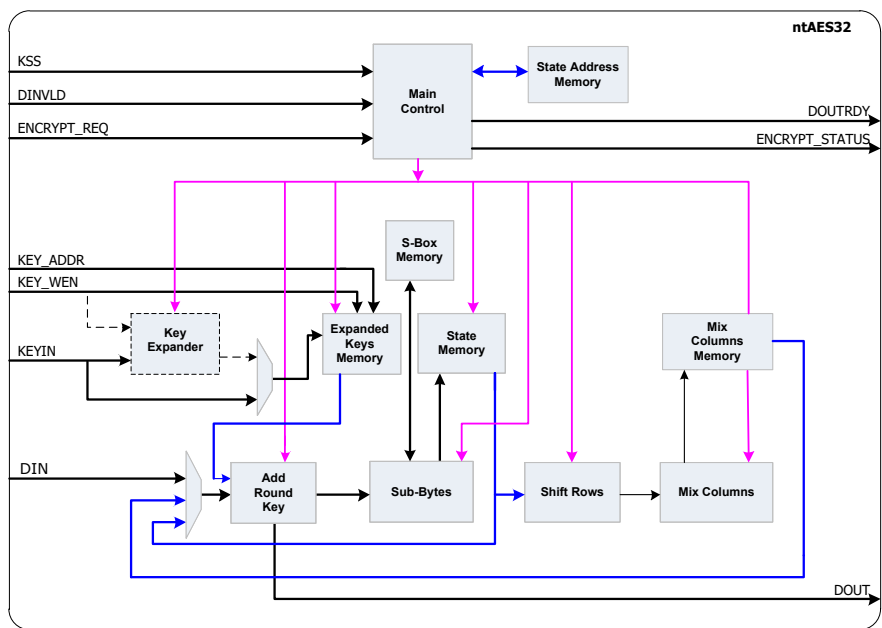| Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Virtex-5	405 CLB Slices / 6 Block RAM	185

Security

ntAES128

AES Ultra High Speed Encryption/Decryption Engine
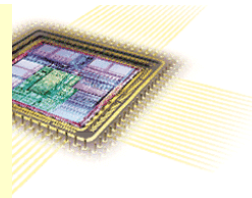
ntAES128 core implements the NIST FIPS-197 Advanced Encryption Standard and can be programmed to either encrypt or decrypt 128-bit blocks of data using a 128-bit, 192-bit or 256-bit key. The ntAES128 has been carefully designed for ultra high throughput applications with optimal logic resources utilization. The encryptor core accepts a 128-bit plaintext input word, and generates a corresponding 128-bit ciphertext output word using a supplied 128, 192, or 256-bit AES key. The decryptor core provides the reverse function, generating plaintext from supplied ciphertext, using the same AES key as was used for encryption. The hardware roundkey expansion logic has been designed as a discrete building block. This allows either to build a complete stand-alone AES solution, or to save logic resources by leaving the key generation process to the user. Alternatively, the roundkey expansion logic can be shared between multiple encryption/decryption cores for optimal silicon area resources utilization. The implementation is very low on latency, high speed with a simple interface for easy integration in SoC applications.

The ntAES128 core can be used in a variety of applications, including:
- Electronic financial transactions.
    eCommerce, Banking, Securities exchange, Point-of-Sale
- Secure communications.
    Storage Area Networks (SAN), Virtual Private Networks (VPN) Video Conferencing, Voice services
- Secure environments.
    Satellite communications, Surveillance systems, Network appliances
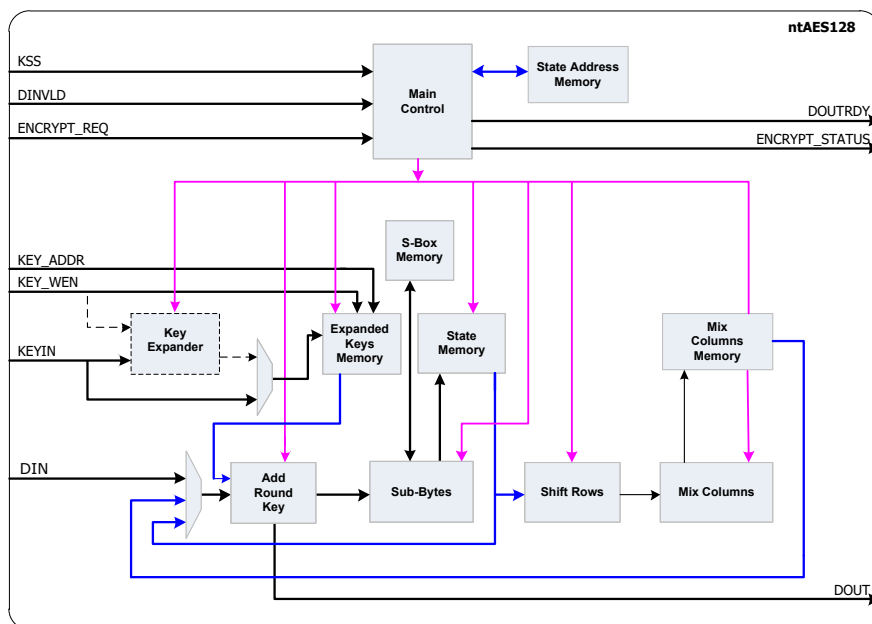- Personal mobile communications.

Key size	Throughput rate (Xilinx Kintex-7)
128 bits	2.25 Gbps
192 bits	1.875 Gbps
256 bits	1.697 Gbps

The ntAES128 core has been core has been implemented to Xilinx FPGA devices.

| Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Kintex-7	570 CLB Slices / 8 Block RAMs	193

ntAES_XTS
XTS Mode AES Processor
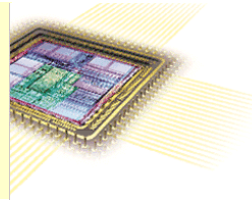
The ntAES_XTS IP Core is fully compliant with AES-XTS algorithm standardized at NIST SP800-38E and IEEE 1619-2007 recommendations targeting disk encryption applications at sector (data unit) addressable level. It is also known as a tweakable block cipher where the encryption process is controlled by the tweak a 128-bit value that is generated from the actual logical position of the data unit on the disk. This way identical data units stored at different places will result in different encrypted data thus addressing copy-and-paste attacks. Each data unit size is at least 128-bits. In addition each data unit size can be either an integral or non-integral number of 128-bit blocks. In case where the data unit size is not divisible with 128 then the ciphertext stealing procedure is used to enable correct encryption of the last block. Due to its highly parameterized and scalable architecture the users can trade off logic resources and performance in order to achieve optimum match with their application requirements. The implementation is low on latency, high speed with a simple interface for easy integration in SoC applications.

The ntAES_XTS core can be used in a variety of applications, including:
- Single SATA 2.0 Hard Disk Drives (up to 3 Gbps throughput rate).
- Single SATA 3.0 SSD (up to 6 Gbps throughput rate).
- USB 3.0 compliant storage.
- Encrypted disk drives.
- SSDs for server arrays (up to 64 Gbps typical throughput rate).
- Encrypted memory sticks.
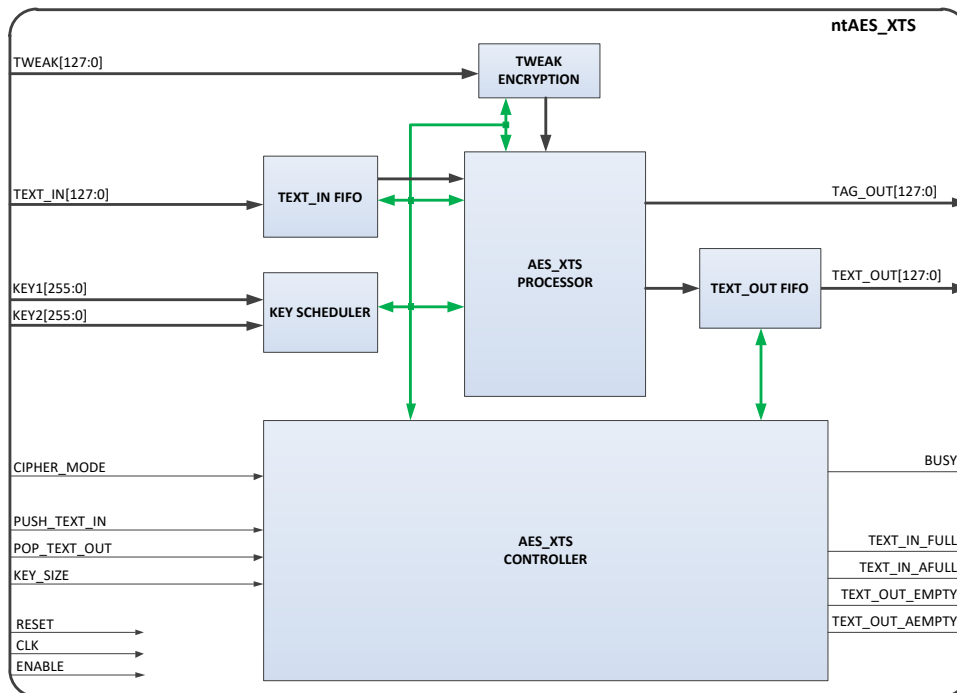
Features
Supports high throughput AES XTS mode for data storage applications.
Compliant with IEEE 1619-2007and NIST SP800-38E recommendations.
Supports 128-bit data-path width.
Supports 128 bit (XTS-256 mode) or 256-bit (XTS-512 mode) key sizes.
Supports cipher stealing mode.
Can be configured either as an encryptor or decryptor mode of operation.
Provides a throughput rate of 16 Gbps at 125 MHz clock rate.
Simple parallel user interface.
Scalable architecture for optimal area/performance trade off.
Fully synchronous design, using single clock.
Silicon proven in ASIC and FPGA technologies for a variety of applications.

ntRC4
RC4 Encryption/Decryption Engine
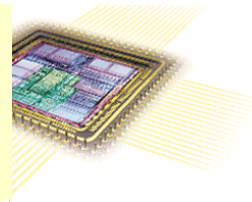
The Noesis Technologies ntRC4 IP core implements the ARC4 stream cipher algorithm. The ntRC4 cipher engine is fully compliant with the wired equivalent privacy (WEP) protocol (part of the IEEE 802.11b wireless LAN security standard) as well as with the IEEE 802.11i (WEP/TKIP). The ntRC4 cipher engine also supports Secure Sockets Layer (SSL) and the companion Transport Layer Security (TLS) standard. It generates the keystream that consists of 8-bit words using a key of length up to 256 bits. The key size length is programmable, the design is fully synchronous with a simple interface that allows seamless integration.

During the key setup phase of the algorithm no input to the core is allowed. The setup phase is completed after 768 clock cycles. KSA_EN is asserted high during the key setup phase. After the START signal is asserted high the encryption/decryption of the plaintext/ciphertext begins and the core produces one encrypted/decrypted byte every three clock cycles. DATA_RDY goes high every 3 clock cycles to indicate that the core is ready to take the next input byte for encryption.
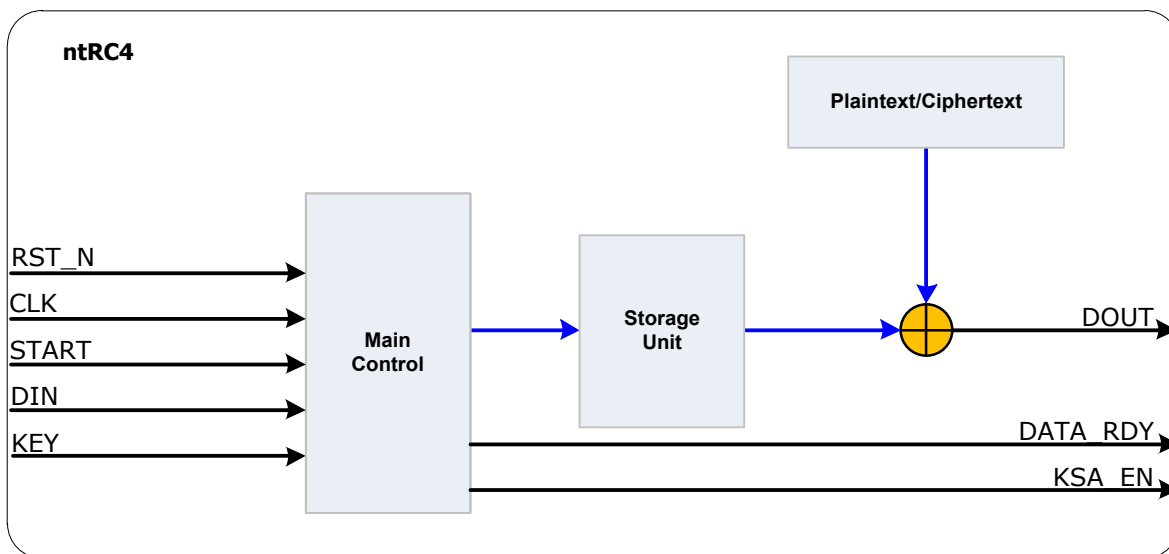
The ntRC4 core has been core has been implemented to Xilinx FPGA devices.

Silicon Vendor	Device	Resources	Throughput Rate	Fmax (MHz)
Xilinx	Virtex-2	134 CLB Slices / 2 Block RAMs	256 Mbps	96
Xilinx	Spartan-6			
Xilinx	Kintex-7			
Altera	Stratix-IV	147 ALUTs / 2 Block RAMs	933 Mbps	350

ntSHA256
SHA 256-bit Hash Generator

An n-bit hash is a map from arbitrary length messages to n-bit hash values. An n-bit cryptographic hash is an n-bit hash which is one-way and collision-resistant. Such functions are important cryptographic primitives used for such things as digital signatures and password protection. Current popular hashes produce hash values of length n = 128 (MD4 and MD5) and n = 160 (SHA-1), and therefore can provide no more than 64 or 80 bits of security, respectively, against collision attacks. Since the goal of the new Advanced Encryption Standard (AES) is to offer, at its three cryptovariable sizes, 128, 192, and 256 bits of security, there is a need for companion hash algorithms which provide similar levels of enhanced security. ntSHA256 IP Core implements SHA-256, or Secure Hash Algorithm-256 which is one of the latest hash functions standardized by the U.S. Federal Government. It is a 256-bit hash and is meant to provide 128 bits of security against collision attacks. The implementation is very low on latency, high speed with a simple interface for easy integration in SoC applications. The ntSHA256 core can used in a variety of applications, including:

- Security applications and protocols (TLS, PGP, SSH, S/MIME, IPsec)
- Authentication of Debian GNU/Linux software packages
- DKIM message signing standard.
- Transaction verification and proof-of-work calculation for several cryptocurrencies (Bitcoin).
- Password protection
- Digital signatures
- Message authentication
- Data integrity check
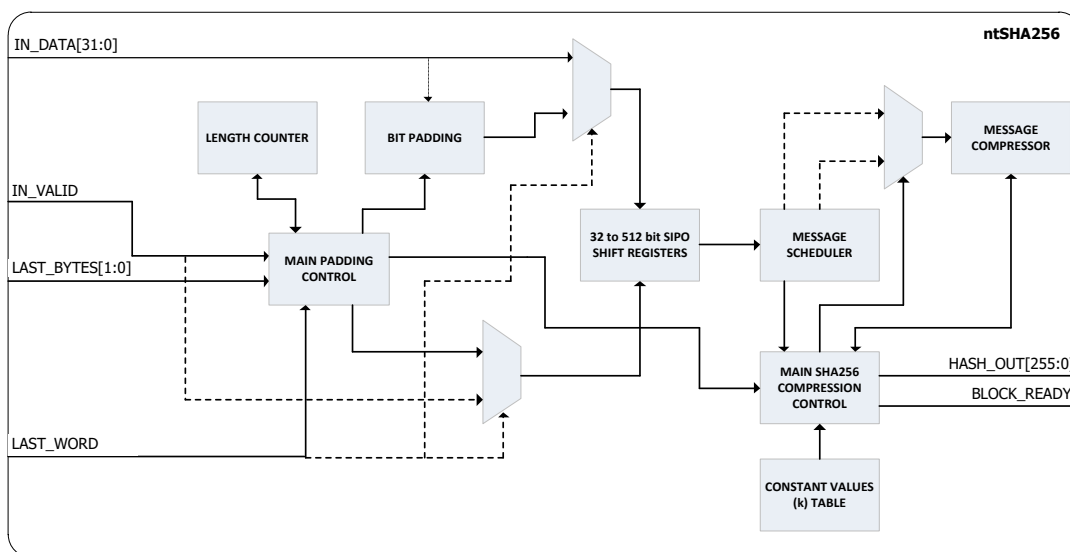
Block size	Throughput rate (Xilinx Spartan-3A)
512 bits	312 Mbps

The ntSHA256 core has been core has been implemented to Xilinx devices.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3A	1577 CLB Slices / 1 Block RAM	50

Telecom DSP Functions

Noesis Technologies offers an extensive library of state-of-the-art signal processing cores used at PHY level of a transmission system.

Discrete Fourier Transforms (DFT) are very common in OFDM based wireless applications as well as in many other telecom applications. The Fast Fourier Transform (FFT) algorithm provides an efficient method for DFT computations in real-time applications. A fully configurable FFT/IFFT processor has been developed that provides SoC designers with a range of high performance FFT cores for various target technologies and application requirements. The ntFFT IP Core employs a revolutionary parameterized architecture where the user can fine tune the level of data-path parallelism in order to achieve the optimum trade-off between silicon resources and throughput rate. A wide range of FFT lengths from 8-point to 8K-points is supported. Support for any power of 2 higher than 8192 is also supported due to fully generic architecture.

A fully configurable soft output demodulator has been designed that receives the equalized complex samples and converts them to a bit stream with soft output information associated with each bit. The probabilistic information is generated based on LLR computations. The core also supports multiple PSK and QAM modulation levels and programmable number of soft bits.

The performance evaluation of a telecom system under the presence of noise using software can be very time consuming. Whereas the noise generation in the analog domain is an easy task, in digital domain the generation of AWGN is a much more complex task. The ntAWGN core has been designed to provide a hardware implementation of an accurate AWGN noise generator that can be used in the efficient performance evaluation of a digital communication system. In addition Noesis Technologies provides customized IP Cores in the areas of channel equalization, channel estimation and synchronization.
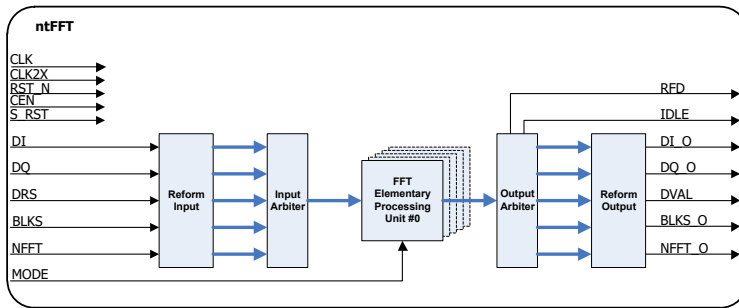
ntFFT
## Fully Configurable FFT/IFFT Radix-2 Processor
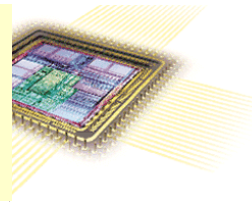
ntFFT core is a fully configurable solution that performs the FFT and IFFT transform. It is on-the-fly programmable in terms of transform size and type. It supports complex input/output and the results are output in normal order. It exhibits a highly parameterizable/scalable design using generic I/O fixed point precision and generic internal calculations precision. The core uses fixed-point 2's complement arithmetic with internal auto scaling to avoid arithmetic overflow and simplify dynamic range management. The ntFFT IP Core employs a revolutionary parameterized architecture where the user can fine tune the level of data-path parallelism in order to achieve the optimum trade-off between silicon resources and throughput rate. The implementation is portable to various silicon technologies, with a simple interface for easy integration in SoC applications.

The modular architecture of the ntFFT processor is shown in the figure below.

The block diagram of the ntFFT elementary processing unit is shown in the figure below.
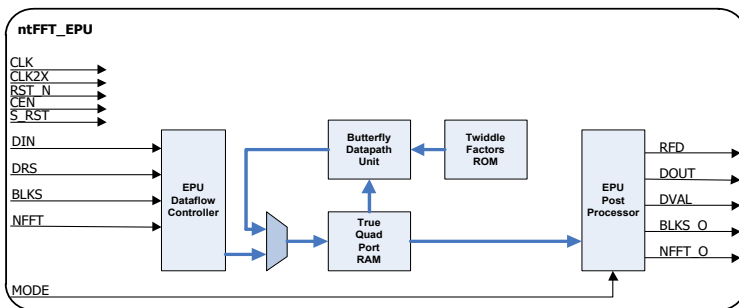
The ntFFT core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details for various FFT sizes are shown in the table below.
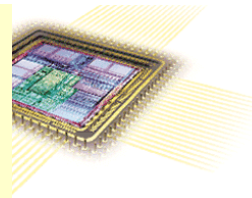
Features
Radix-2 Fast Fourier Transform processor IP Core.
Same IP core may be used to compute both FFT and IFFT transforms without any complexity overhead.
Highly parameterizable/scalable design using generic I/O fixed point precision and generic internal calculations precision.
Bit true Matlab script model is provided to aid core fixed point precision configuration for any target application.
Tested against Matlab FFT and IFFT functions the ntFFT core produces fixed point numerical results with mean absolute error in range of 1e-4. The core may be parameterized for greater internal fixed point precision to lower the mean absolute error further.
Final fixed point scaling to avoid precision loss is performed internally.
Highly programmable design supporting all power of 2 FFT/IFFT transforms in range [8, ...,MAX_NFFT], where MAX_NFFT=[8,...,8192]. Support for any power of 2 higher than 8192 is also possible.
Parameterized architectural parallelism level to meet any target application by tuning an efficient trade-off between utilized resources and maximum throughput rate.
Overclocked main memory at 2x rate to achieve minimum memory resources utilization.
Simple yet robust interface for optimum and efficient data flow control.
Optional AXI4-Stream protocol interface support.
Fully synchronous design.
Silicon proven in ASIC and Xilinx FPGA implementation technologies for a variety of applications.

| Silicon Vendor	Device	FFT size	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	128	547 CLB Slices / 4 Block RAMs / 8 DSP48 Blocks	206
Xilinx	Kintex 7 XC7K410T-2	512	716 CLB Slices / 5 Block RAMs / 8 DSP48 Blocks	190
Xilinx	Kintex 7 XC7K410T-2	1024	652 CLB Slices / 6 Block RAMs / 12 DSP48 Blocks	171
Xilinx	Kintex 7 XC7K410T-2	2048	765 CLB Slices / 16 Block RAMs / 16 DSP48 Blocks	170
Xilinx	Kintex 7 XC7K410T-2	4096	745 CLB Slices/ 16 Block RAMs / 16 DSP48 Blocks	170
Xilinx	Kintex 7 XC7K410T-2	8192	836 CLB Slices/ 33 Block RAMs / 16 DSP48 Blocks	166

ntCH_EST
Programmable OFDM Channel Estimator

The wideband OFDM signal suffers from frequency selective fading. Therefore it is necessary to identify and invert the discrete transfer function of the channel. The accurate channel estimation is achieved with the exploitation of known reference signals and pilots into the OFDM frame. The ntCH_EST core uses the pilots to determine the channel impulse response in the frequency domain. Channel estimation is performed on a block-per-block basis, where one block is composed of a programmable number of OFDM symbols. The pilot allocation and the block size is fully programmable. The ntCH_EST implements estimation formulas based on Linear Least Squares (LS) and 1D linear interpolation algorithms for optimum trade-off between complexity and accuracy

Specifically the channel estimation performs operations on the specific block which are described by the following algorithmic steps:

- Computation of the expected pilot positions and the expected pilot modulation.
- Isolation of the pilot subcarriers from the incoming signal.
- Averaging of the pilot values for the selected channel estimation block to achieve better estimation results.
- Applying the channel estimation formula and calculating the discrete frequency transfer function value for each pilot.
- Interpolating, using linear interpolation techniques, the estimated values in the frequency domain to extract the transfer function for the data subcarriers.

The ntCH_EST supports programmable pilot patterns and programmable OFDM frame size. It is a fully synchronous design, using single clock. It is silicon proven in ASIC and FPGA technologies for a variety of applications.
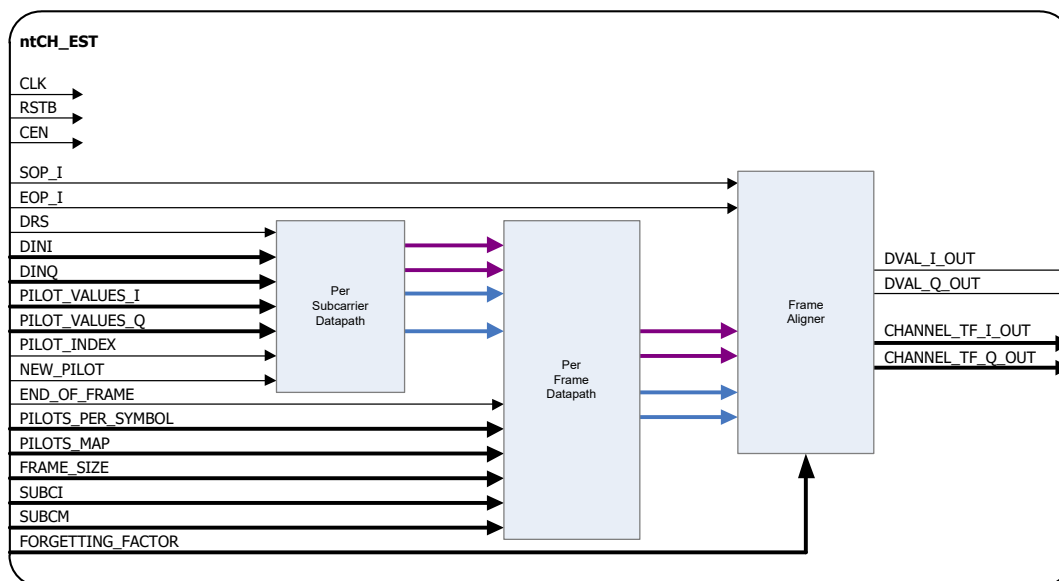
The ntCH_EST core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details for 256 subcarriers per OFDM symbol configuration are shown in the table.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	256 subcarriers / OFDM symbol	2495 CLB Slices / 8 Block RAMs / 44 DSP48 Blocks	100

ntSOD
Fully Configurable BPSK, QPSK, QAM Soft Output Demapper

Noesis Technologies ntSOD Soft Output Demapper is a structural element of any modern telecom system. The receiver extracts the phase and magnitude of the carrier signal. Subsequently a decision must be taken on the actual transmitted bits. Due to channel noisy conditions, the received signal has been distorted and there are positional errors on the constellation points. The ntSOD Soft Output Demapper IP Core implements the LLR (Log Likelihood Ratio) algorithm to convert the received distorted modulated signal from its complex I, Q form to a bit stream. It identifies the actual transmitted symbol bits and assigns to each bit a level of confidence in the format of a soft value. It supports various modulation levels such as BPSK, QPSK, 16 QAM and 64 QAM. This soft-bit information can be subsequently used during ECC decoding process by a soft-input ECC decoder such as Viterbi Decoder. Soft decision ECC decoding can provide a coding gain of 2 dB for 3 soft-bits per encoded bit or 2.2 dB for 4 soft bits per encoded bit when compared with hard decision ECC decoding. The soft-bit information can be configured in sign-magnitude or 2's complement format. The number of soft bits per symbol are parameterized as well as the supported modulation levels. It is a fully synchronous design, using a single clock.
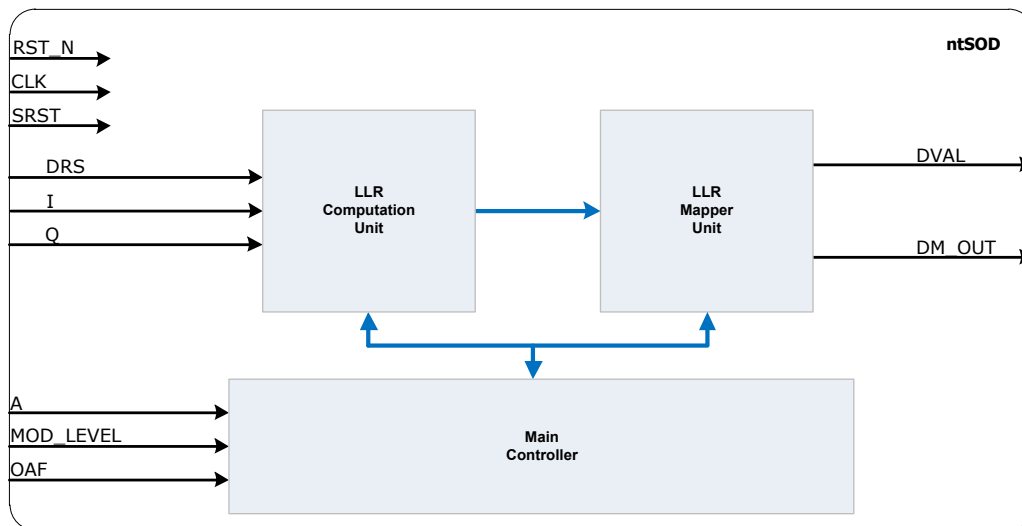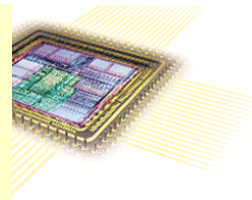

The ntSOD core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	BPSK up to QAM64	352 CLB Slices / 6 DSP48 Blocks	130

ntAWGN
AWGN Channel Emulator

The performance evaluation of a telecom system under the presence of noise using software can be very time consuming. Whereas the noise generation in the analog domain is an easy task, in digital domain the generation of AWGN is a much more complex task. The ntAWGN core provides an innovative all digital hardware implementation of a highly accurate AWGN noise generator that can be used in the efficient performance evaluation of a digital communication system. The core generates AWGN with the following characteristics:

- Based on the Box-Muller algorithm and the Central Limit Theorem.
- Random distribution in the range of [-4σ...4σ], where σ is the standard deviation with probability density function (PDF) deviation less than 1 % from the Gaussian.
- Noise gain precision 14 bits allows accurate resolution in the range of [0,30] db SNR.
- Periodicity up to 2^{60} samples.
- Noise samples generated from Box-Muller engine are 10 bits wide (NOISE_WIDTH), with 4 bits of integer and 6 bits (NOISE_FRAC) of fractional part.
- Generics allow modification of arithmetic precision, number of accumulations, LFSRs initialization.

Bit errors are generated by adding a white gaussian noise variable to the input bit stream. The number of bit errors and therefore the noise level is controlled by adjusting the standard deviation of the AWGN and/or the input signal amplitude. The ntAWGN core is comprised of configurable number of independent white gaussian noise generators that are used to add noise to the incoming signal represented by 10-bit per sample precision. The following figures demonstrate the AWGN probability density function and BER vs SNR correlation between theoretical (red line) and real-time measurements (blue line).
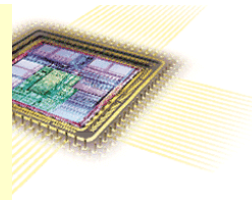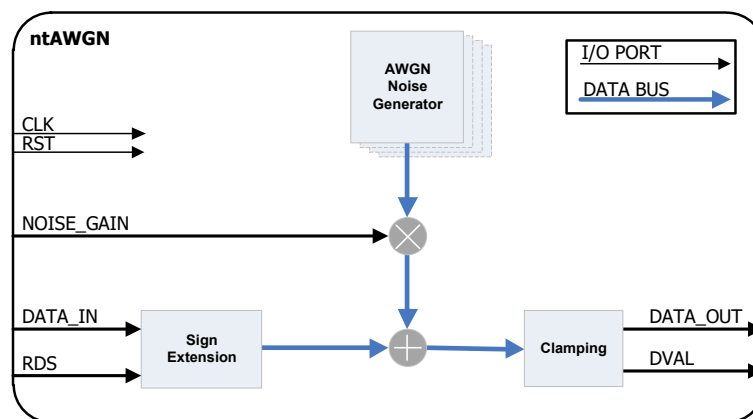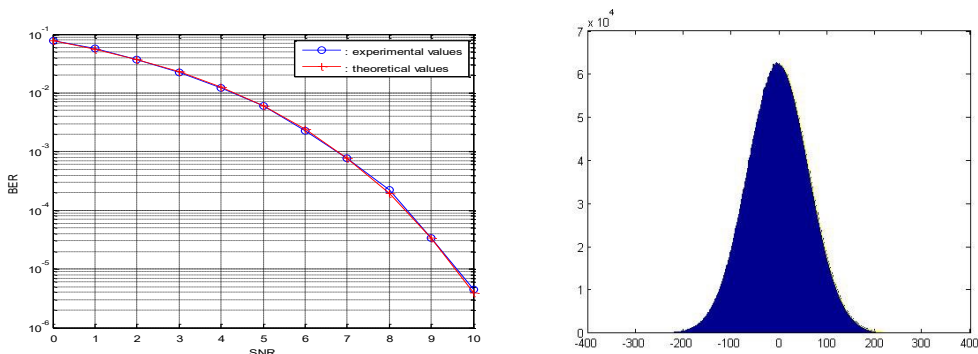

The ntAWGN core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	265 CLB Slices / 4 Block RAMs / 8 DSP48 Blocks	360

ntSYNC
Synchronization Unit

Noesis Technologies ntSYNC is a fully programmable component used to achieve time and frequency synchronization in OFDM technology physical layer implementations. It interfaces directly with the physical layer's front-end (line/RF) interface and using a cross-correlation proprietary algorithm to find the starting point of a received data frame. The generic design approach as well as a number of pre-processed optimizations allow for integration to any OFDM compliant physical layer.

The front-end interface feeds the ntSYNC with received signal above a certain power level via the DINI/DINQ ports and flags the signal as valid (DRS). The signal is buffered temporarily in Raw Buffer until raw synchronization phase takes place. The raw synchronization algorithm searches for received power level equal to the a-priori known preamble power levels. The power levels can be programmed via the PWR_THRES input port.

Once raw synchronization is achieved the estimated location of the preamble is decided and the coarse synchronization phase begins. Coarse synchronization searches the preamble estimated location more closely with correlative metrics, approximates the received frames starting point and calculates the channels frequency offset shift on the received signal. The raw buffer discards data before the approximated synchronization point and propagates the rest of the signal to the frequency offset compensation unit. The results are being stored in the fine buffer. Additionally the estimated preamble data are being isolated.

The isolated preamble is correlated against the known generated preamble by the fine synchronization process. Fine synchronization decides the final synchronization point and discards all previous data in the fine buffer. The programmable OFDM control parameters such as cyclic prefix size (CP), sub-channelization operation (SUBCI), the size of the OFDM symbol (OFDM_SIZE) and the number of OFDM symbols that are expected to be included in the received frame (OFDM_NUM) are required in order to decide the end of the received frame, once the synchronization point has been calculated.
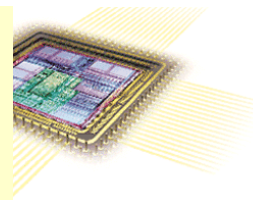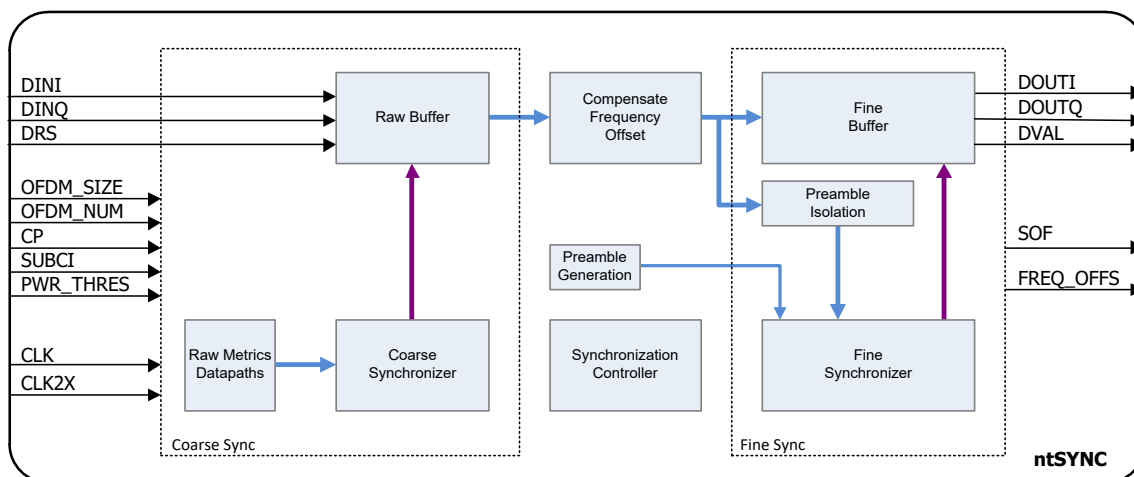

The ntSYNC core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	5314 CLB Slices / 27 Block RAMs / 96 DSP48 Blocks	86 (CLK) / 102 (CLKX2)

Communication Protocols & Networking
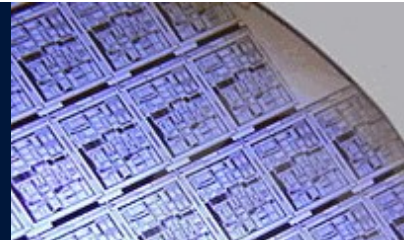
Noesis Technologies has developed a broad library of channelized E1/T1, E2 & E3 framer/de-framer IP Core solutions that can be used in a variety of time division multiplexing (TDM) applications including wireless base transceivers, digital subscriber line access multiplexers, voice gateways, private branch exchanges (PBXs), optical network units (ONUs), integrated access devices (IADs), routers and test equipment. The ntE1_G704 and ntT1_G704 cores are fully compliant with Dallas Semiconductor DS2186/DS2187 transmit/receive line interface units. In addition Noesis Technologies, offers the ntHDLC core that implements a single-channel controller for the High-Level Data Link Control (HDLC) protocol that can be used in public networks employing the X.25 communications protocol, xDSL transport, frame relay and ISDN applications.

ntE1_G704

E1 Framer/Deframer—ITU-T G704 compliant
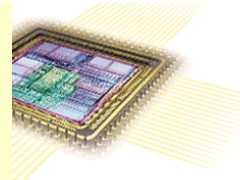
Noesis Technologies ntE1_G704 Framer/Deframer is designed for E1 networks and is compliant with ITU recommendations G.704, G.706, G.732, G.775 and O.163. The core provides all the necessary data formatting transforms for transmission over an E1 carrier. E1 is one of the two most widely used TDM (time division multiplexing) carriers incorporating 32 channels, each with a bandwidth of 64 kbps providing a total bit rate of 2048 kbps. The ntE1_G704 IP core provides a flexible interface supporting hardware and microprocessor modes. Specifically the core can be connected to a host system either through an 8-bit parallel microprocessor interface (HP mode) or through a set of I/O ports (HW mode). When in HP mode, the microprocessor configures and monitors the functionality of the core through a rich set of registers. When in HW mode, the core is directly controlled and monitored through a set of dedicated ports and no microprocessor control is necessary. At the transmit side, the framer generates framing patterns, CRC4 bits, formats outgoing and signaling data, generates alarms and clock outputs for data conditioning and decoding. At the receive side, the deframer establishes frame / multiframe synchronization, extracts data, signalling and alarm flags. It provides information like frame, multiframe alignment, calculates CRC4, counts CRC4 errors and performs A-bit processing.
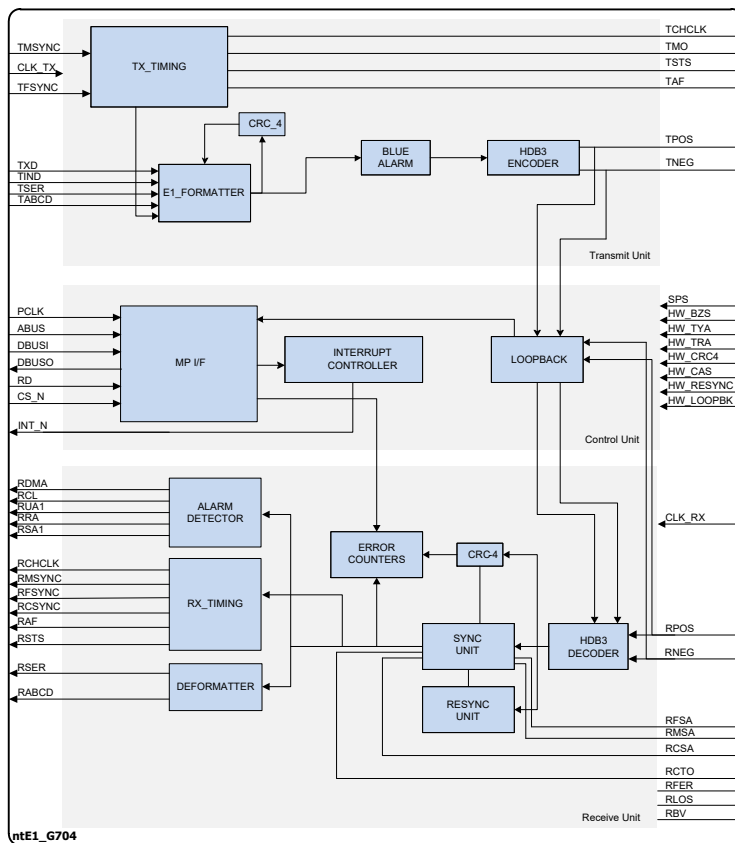
Features

Features
E1 framer/deframer compliant to G.704, G.706, G.732, G.775 and O.163 CCITT recommendations.
Supports CAS and CCS signaling standards.
Supports CRC4 based framing standards.
User configurable receive and transmit control.
Supports 8-bit parallel microprocessor interface for device configuration and control in host processor mode.
Hardware control mode requires no host processor; ideal for stand-alone applications.
Supports HDB3 line coding.
Supports loop-back mode.
Alarm generation, alarm detection and error logging.
Compatible with Dallas DS2186 transmit line interface and DS2187 receive line interface.
Fully synchronous design.
Silicon proven in ASIC and Xilinx FPGA implementation technologies for a variety of applications.

The ntE1_G704 core has been core has been implemented to Xilinx and Altera FPGA devices as well as TSMC ASIC libraries.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3	1027 CLB Slices	102 (clk_tx)/ 107 (clk_rx)/ 207 (pclk)
Altera	Stratix-III	887 ALUTs	230 (clk_tx) / 244 (clk_rx) / 480 (pclk)
TSMC	180 nm	9200 gates	400 (clk_tx) / 432 (clk_rx) / 770 (pclk)

ntT1_G704
T1 Framer/Deframer—ITU-T G704 compliant

Noesis Technologies ntT1_G704 Framer/Deframer is designed for T1 networks and is compliant with ITU recommendations G.704, G.706, G.732, G.775 and O.163. The core provides all the necessary data formatting transforms for transmission over an T1 carrier. T1 is one of the two most widely used TDM (time division multiplexing) carriers incorporating 24 channels, each with a bandwidth of 64 kbps providing a total payload bit rate of 1536 kbps. The ntT1_G704 IP core provides a flexible interface supporting hardware and microprocessor modes. Specifically the core can be connected to a host system either through an 8-bit parallel microprocessor interface (HP mode) or through a set of I/O ports (HW mode). When in HP mode, the microprocessor configures and monitors the functionality of the core through a rich set of registers. When in HW mode, the core is directly controlled and monitored through a set of dedicated ports and no microprocessor control is necessary. At the transmit side, the framer generates framing patterns, CRC6 bits, formats outgoing and signaling data, generates alarms and clock outputs for data conditioning and decoding. At the receive side, the deframer establishes frame / multiframe synchronization, extracts data, signalling and alarm flags. It provides information like frame, multiframe alignment, calculates CRC6 and counts CRC6 errors.
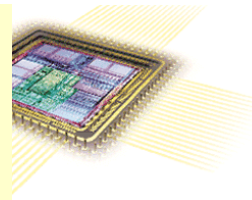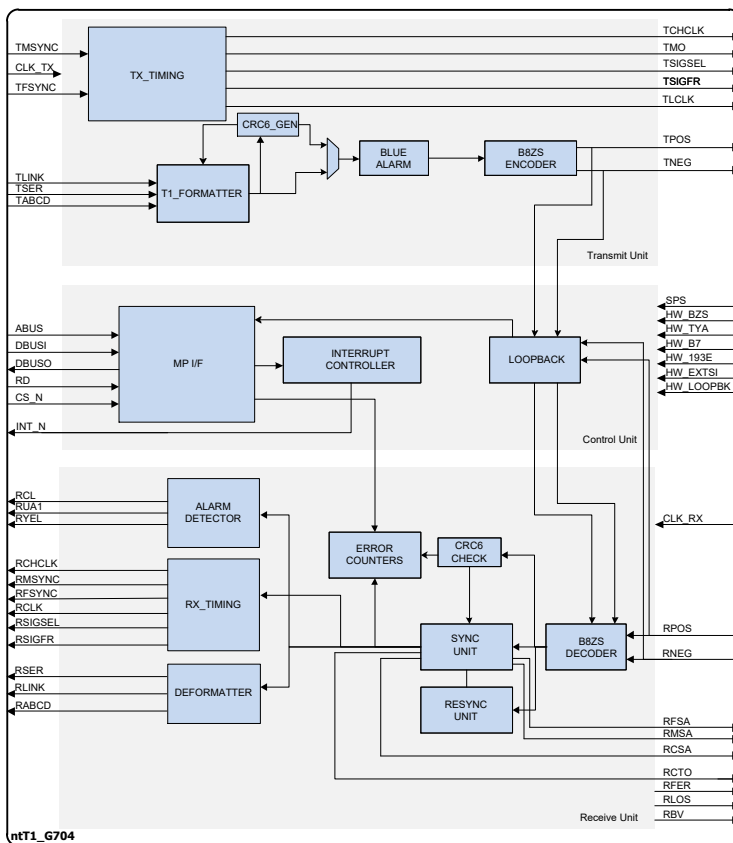

Features

Features
T1 framer/deframer compliant to G.704, G.706, G.732, G.775 and O.163 CCITT recommendations.

Supports common framing standards.
   - 12 frames/superframe '193S'
   - 24 frames/superframe '193E'

Supports CRC6/non-CRC6 cases.
Three zero suppression modes: B7 stuffing, B8ZS and transparent.
User configurable receive and transmit control.
Supports 8-bit parallel microprocessor interface for device configuration and control in host processor mode.
Hardware control mode requires no host processor; ideal for stand-alone applications.
Supports loop-back mode.
Alarm generation, alarm detection and error logging.
Compatible with Dallas DS2186 transmit line interface and DS2187 receive line interface.
Fully synchronous design.
Silicon proven in ASIC and Xilinx FPGA implementation technologies for a variety of applications.

The ntT1_G704 core has been core has been implemented in Xilinx FPGA devices as well as in TSMC ASIC libraries.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3	1027 CLB Slices / 12 Block RAMs	101 (clk_tx)/ 100 (clk_rx)
TSMC	180 nm	7050 gates NAND2 equivalent / 4632 RAM bits	432 (clk_tx) / 350 (clk_rx)

ntE2_E3

E2 & E3 Framer/Deframer—ITU-T G742/G751 compliant

Noesis Technologies ntE2_E3 Framer/Deframer is designed for E2/E3 networks and supports all requirements of ITU recommendations G.742, G.751 and G.775. The core provides all the necessary data formatting transforms for transmission over E2/E3 networks. The device can by controlled through a simple set of dedicated ports, allowing robust operation. One ntE2_E3 core instance can operate either as an E1/E2 (2.048/8.448 Mbps) rate Multiplexer/Demultiplexer, or as an E2/E3 (8.448/34.368 Mbps) rate Multiplexer/Demultiplexer. In addition five ntE2_E3 cores can be instantiated to operate as an E1/E3 (2.048/34.368 Mbps) rate Multiplexer/Demultiplexer. The transmit side of the framer generates framing patterns, transmits the alarm and the national bits, interleaves the four tributaries into the high level data stream, calculates the justification mechanism status and the nature of the stuffing bits available, as well as generates alarms, status bits, and clock outputs. The receive side establishes frame synchronization, extracts the interleaved data, the alarm and national bits as well as the auxiliary channels data, monitors for error conditions and generates alarm flags, data valid bits, status bits and clock outputs. The HDB3 codecs can be either used or bypassed, on both transmit and receive sides, depending on the application. Finally both local and remote loopback features are available.
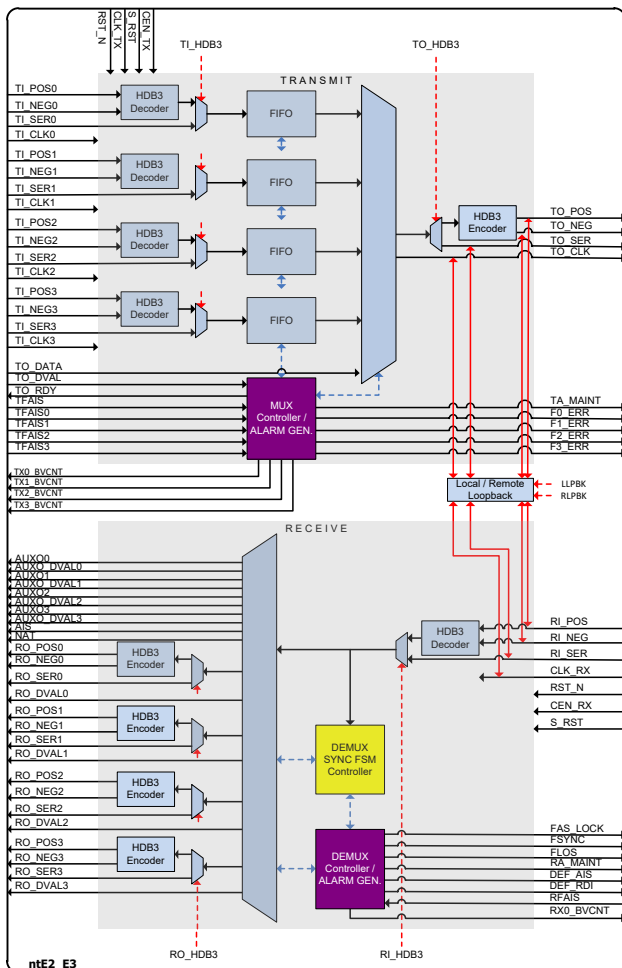
Features

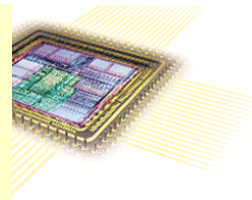| Features |
| --- |
| E2/E3 framer/deframer compliant to G.742, G.751, G.775 ITU-T standards. |
| Performs four E1 to one E2 or four E2 to one E3 multiplexing and vice-versa demultiplexing. |
| Five ntE2_E3 cascaded cores implement a sixteen E1 to one E3 Multilpexer/Demultilpexer. |
| Optional HDB3 Line Codecs one both Receive and Transmit sides. |
| Local and Remote Loop-back modes. |
| Configurable Frame Alignment Signal. |
| User access to the Alarm bit and the National bit. |
| User access to four low speed Auxiliary Channels, one per multiplexed tributary, available via unused Stuffing bits. |
| Alarm generation, alarm detection and error logging. |
| Fully synchronous and parametric design. |
| Silicon proven in ASIC and FPGA technologies for a variety of applications. |

The ntE2_E3 core has been core has been implemented in Xilinx FPGA devices.

Silicon Vendor	Device	Configuration	Resources	Fmax (MHz)
Xilinx	Virtex-6	E2 mode (2Mbps/8 Mbps)	436 CLB Slices	170 (clk_tx)/ 250 (clk_rx)
Xilinx	Virtex-6	E3 mode (8Mbps/34 Mbps)	406 CLB Slices	186 (clk_tx) / 205 (clk_rx)

ntHDLC

High Level Data Link Controller

Noesis Technologies ntHDLC single channel High-Level Data Link Controller (HDLC) is a full-duplex transceiver with independent transmit and receive units for synchronous framing bit-level HDLC protocol operations. The ntHDLC can handle interframe and delimiting flags, frame check sequence based on CCITT CRC16/CRC32 polynomial, normal or transparent transmission modes, abort generation and detection. The system interface is very flexible and can be adapted towards FIFO, uP, or DMA controllers. The transmit and receive units and their associated control and status logic are independent. This partitioning strategy enables the Tx and Rx units to be instantiated in different place and/or level of the design hierarchy. Each unit (Tx, Rx and back-end interface) has its own clock domain with synchronous clock enable. Communication between the various clock domains is achieved via synchronization logic blocks.
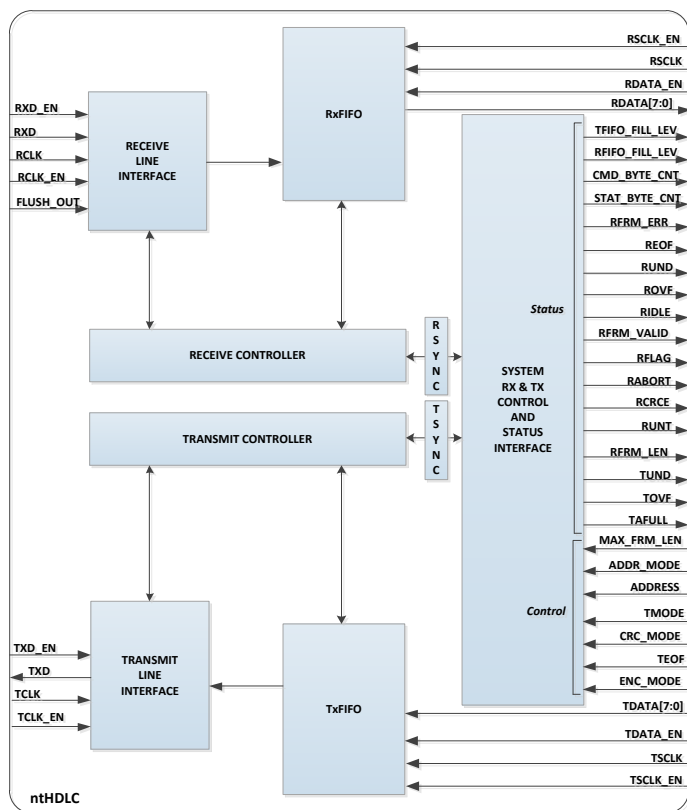
Features

| Features |
|---|
| Single port synchronous serial line interface. |
| Flag/Abort Generation/Detection. |
| Zero Insertion/Deletion. |
| Non-octet alignment detection. |
| CCITT CRC-16 Generation and Checking. |
| NRZ/NRZI encoding/decoding. |
| Transparent mode support. |
| Receive FIFO overrun detection. |
| Transmit FIFO underrun detection. |
| Frame status and frame length indicators. |
| Runt frame detection. |
| Seperate clocks for Tx and RX interfaces. |
| Supports flag in interframe-time fill. |
| 8-bit parallel back-end interface. |
| Fully synchronous design. |
| Silicon proven in ASIC and Xilinx FPGA implementation technologies for a variety of applications. |

The ntHDLC core has been core has been implemented to Xilinx and Altera FPGA devices as well as TSMC ASIC libraries.

| Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Spartan-3	460 CLB Slices / 10 Block RAMs	80 (tclk)/ 126 (rclk)/ 140 (tsclk)/124 (rsclk)
Xilinx	Virtex-5	200 CLB Slices / 10 Block RAMs	130 (tclk)/ 230 (rclk)/ 310 (tsclk)/313 (rsclk)
Altera	Stratix-III	600 ALUTs / 10 M9K RAM Blocks	72 (tclk)/ 139 (rclk)/ 184 (tsclk)/133 (rsclk)
TSMC	180 nm	5800 gates NAND2 equivalent / 74 K RAM bits	340 (tclk)/ 400 (rclk)/ 330 (tsclk)/340 (rsclk)

Baseband PHYs
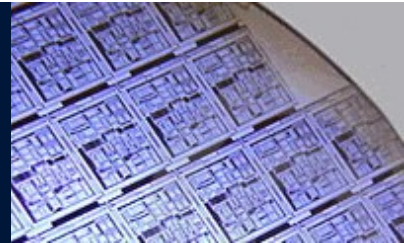
OFDM transmission technology is spectrally efficient and very robust to harsh wireless channel environments. It is widely applied in wireless communication systems providing high rate transmission capability coupled with high bandwidth efficiency as well as robustness to multi-path fading and multi-path delay. Its frequency selectivity feature allows the users to disable certain OFDM subcarriers to prevent interfering with the others and makes the technology extremely robust against frequency selective fading transmission conditions.

Noesis Technologies ntOFDM_BBP is a custom baseband processor, which implements the physical layer of an OFDM, time division duplexing (TDD) system. The baseband processor includes both transmission and reception bit-level and symbol-level processing chains including a sophisticated synchronization unit. The host interface is based on AXI-stream protocol. This custom system implements a subset of 802.16d standard functional options/features and is highly configurable via the integrated register-file. An RF interface module is also included, compatible with Analog Devices AD9361 RF transceiver.

ntOFDM_BBP

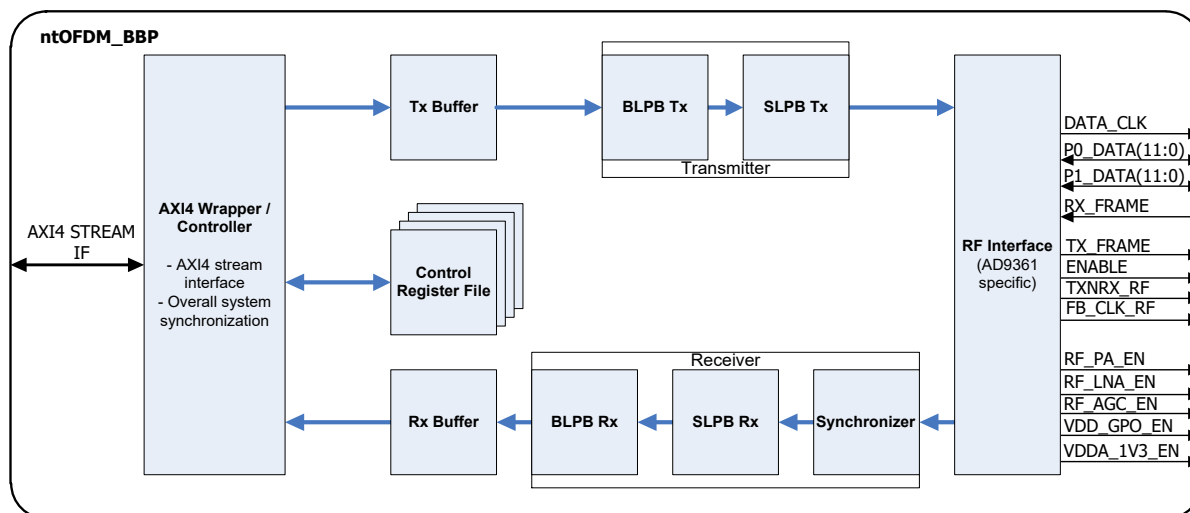## Multi-Purpose OFDM Baseband Processor
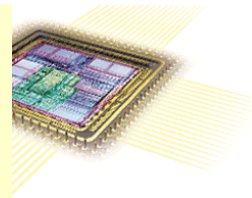
Noesis Technologies ntOFDM_BBP is a custom baseband processor, which implements the physical layer of an OFDM, time division duplexing (TDD) system. The baseband processor includes both transmission and reception bit-level and symbol-level processing chains including a sophisticated synchronization unit. The host interface is based on an AXI4 stream protocol. This high performance OFDM transmission system is fully compliant with 802.16d (WiMAX) standard and is fully configurable via the integrated register file. An RF interface module is also included, compatible with Analog Devices AD9361 RF transceiver. Other RF interfaces can be supported.

The Bit-level processing block (BLPB) transmission chain implements the following functional units : randomization, FEC encoding, interleaving and symbol mapping. In BLPB reception chain the following operations are implemented: soft symbol demapping, de-interleaving, FEC decoding and de-randomization. The FEC module implements a powerful error correction scheme based on a concatenation of Reed Solomon—Viterbi algorithms.

Features
Customized transmit and receive physical layer chains.
Fully synchronous design enabling high throughput TDD operation.
BLPB and SLPB processing blocks.
Implements a sophisticated synchronization algorithm to efficiently detect and isolate received modulated payload information.
Configurable as either downlink (DL) baseband station or uplink (UL) baseband station.
Configurable data randomization, modulation level and code rate.
Decoding algorithm achieves competitive performance results with the minimum possible test patterns and decoding iterations.
Host interface based on AXI4 stream protocol.
RF interfacing supporting Analog Devices AD9361 RF Transceiver.
Synchronous single clock design.
Silicon proven in ASIC and Xilinx FPGA implementation technologies.

The Symbol-level processing block (SLPB) transmission chain implements the following functional units: OFDM symbol transmitter, IFFT, CP insertion. In reception chain the SLPB module is preceded by the synchronization unit, which is searching for known preamble values in order to locate the start an incoming WiMAX sub-frame. Once the sub-frame is located, frequency offset compensation is applied and received information is propagated down to SLPB reception chain. In SLPB reception chain the following operations take place: CP removal, FFT, OFDM symbol receiver, channel estimation, phase offset compensation and channel equalization.


The ntOFDM_BBP core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details are shown in the following table.

Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	24763 CLB Slices / 128 Block RAMs / 294 DSP48 Blocks	100

ntOFDM_BBP
OFDM Baseband Processor
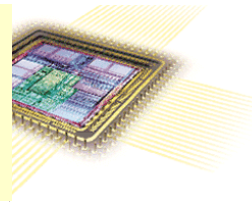
The ntOFDM_BBP advanced DSP algorithms eliminate channel white noise effects as well as frequency and phase offset distortions. BER vs Eb/No performance curves are shown in the figure below for an AWGN channel, for BPSK, QPSK and QAM16 modulation levels as well as for 1/2 and 3/4 coding rates.
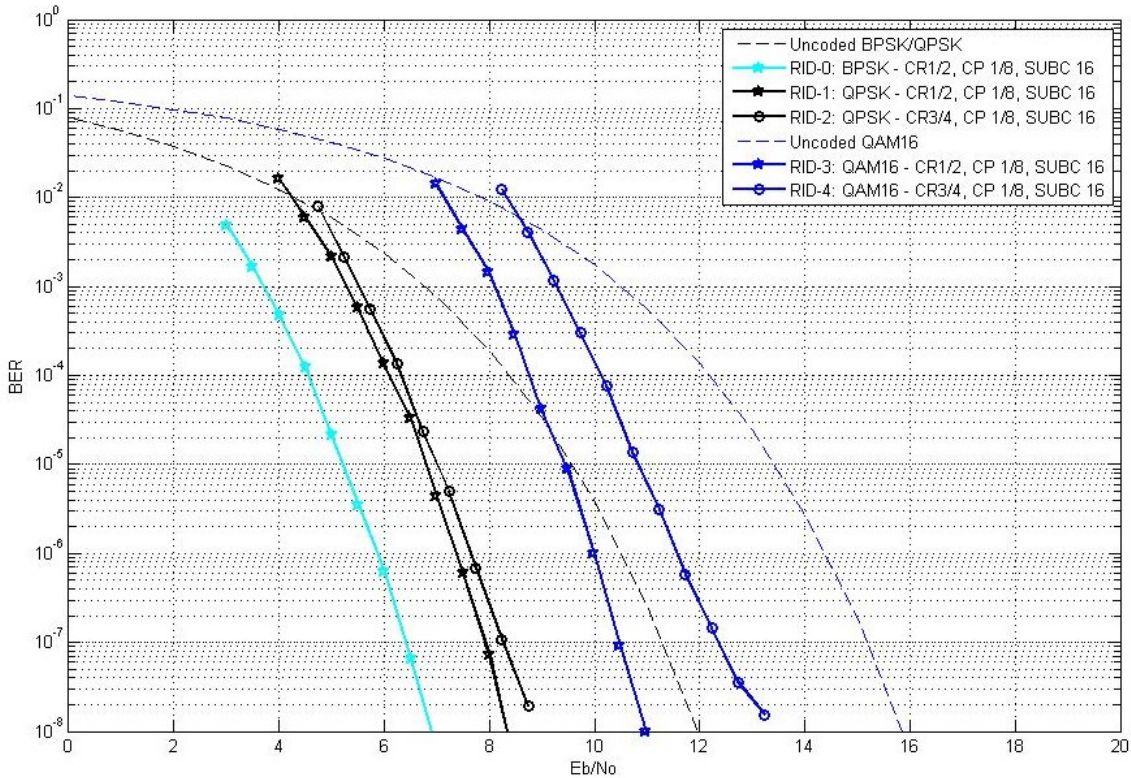
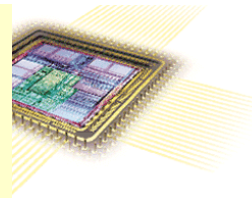The following table presents for various operation modes the achievable throughput rates for 16-sub-channels (full sub-channelization mode), 1/8 cyclic prefix, TDD mode and 50 MHz system clock frequency.

Operation mode	Throughput rate (Mbps)
BPSK—Code rate 1/2	12
QPSK—Code rate 1/2	25
QPSK—Code rate 3/4	33
QAM16—Code rate 1/2	31
QAM16—Code rate 3/4	38

ntGhn_BBP
Home PLC Baseband Processor (under dev)

Noesis Technologies ntGhn_BBP is a fully ITU-T G.9960/G.9964 compliant baseband processor, featuring a highly innovative configurable architecture that enables unprecedented power-line throughput rates. It can be used in a variety of applications including smart grid and home automation, high bandwidth home networking, IPTV infrastructure, consumer electronics.
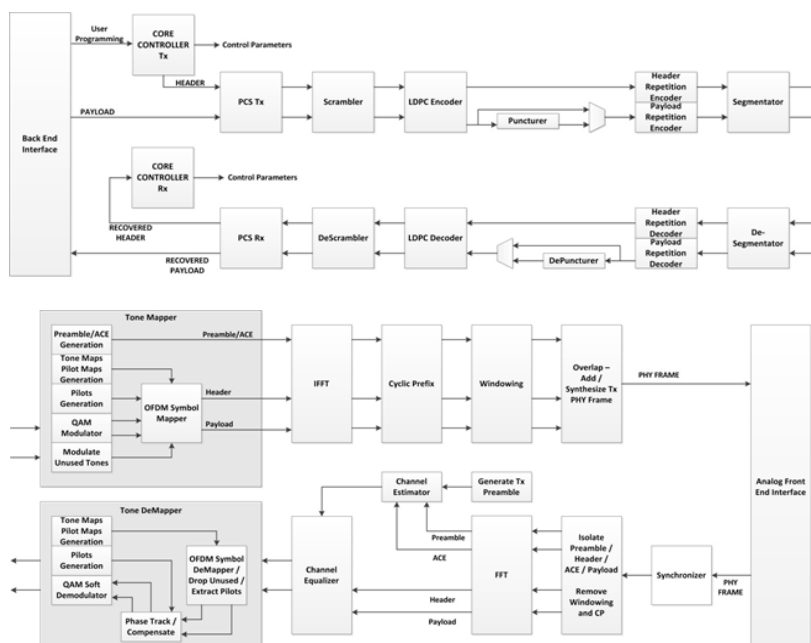
The ntGhn_BBP IP core main functional blocks are the transmitter, the receiver, the back-end interface module and the analog front end (AFE) interface. The back-end interface may be optionally wrapped with an AMBA AXI 4 Stream controller. Both the transmitter and the receiver include a local controller and a register file. Via the back-end interface the user may write or read data and configure the cores to operate with a specific functional mode of operation (control profile).

The Transmitter PCS component encodes the active register file control information to a Header data block and computes all derivative control configuration parameters required for Payload transmission. As per G.9960 both Header and Payload data blocks are scrambled, LDPC encoded, (optionally) repetition encoded and finally partitioned to Symbol Frames in the PMA component of the Transmitter. The PMD component of the Transmitter modulates and tone maps the symbol frame data segments to OFDM symbols. The OFDM symbols are then processed by the constellation scrambler and the IFFT, which transforms the signal from the frequency to the time domain. The time domain signal is cyclic prefix extended, windowed and finally overlapped to form the final PHY Frame. The formed PHY Frame is provided to the AFE transmit interface, which synchronizes the symbols to the front end clock rate.

The AFE receive interface passes data from the channel to the Synchronizer unit. The Synchronizer is responsible of sensing the channel conditions and establishing a synchronization point in time, in other words it locates the start of the PHY Frame by searching for the Preamble symbols energy levels. After synchronization the Receiver PMD component removes the effects of windowing and the cyclic prefix, applies the FFT transform to cross the signal back to the frequency domain and thoroughly estimates the channel conditions in terms of phase and frequency offsets. Then the Receiver equalizes and compensates the received OFDM symbols accordingly. Once the Header part of the PHY Frame is identified and equalized, it is being demodulated by the tone demapper and passed to the Receiver PMA component for the optional repetition decoding, LDPC decoding and de-scrambling. A CRC check is performed on the received Header information in the Receiver PCS component and if successful, the Payload reception mode of operation is configured. All derivative control configuration parameters required for payload reception are computed by the Receiver PCS component. Then the Payload OFDM symbols are propagated through the remaining Receiver PMD and PMA components according to the selected control configuration. Finally the decoded Payload data blocks are being returned to the User via the back-end interface.

Features
G.hn physical layer (PHY) baseband processor is compliant to ITU-T G.9960 and ITU-T G.9964 standards
Supports telephone line, power line and coaxial (baseband and RF) bandplans
Supports MSG, ACK, RTS and CTS Frame Types and has infrastructure to easily implement all other G.9960 defined frame types
Highly programmable core supports all standard defined configuration profiles and additionally enables construction of custom configuration profiles
Expandable and highly configurable architecture via a set of generic values with several degrees of parallelism for optimum design approach
Architecture programmability and expandability provide a wide range of trade-offs between core area utilization and information throughput rates
Gb/s information throughput rates can be achieved by expanding the architectural parallelism and selecting high data rate modes of operation in optimum channel conditions, for operating frequency as low as 100MHz
Operating frequency is target technology dependent and raises achievable information throughput proportionally
High speed design approach accomplishes operating frequency greater than 100MHz in medium speed grade FPGA Kintex7 prototyping boards
Distributed control pipelines and flow controller mechanism to achieve minimum component flush out times and "variable on-the-fly" configurability
Simple and robust core back end interface enables high throughput data flow
Optional "AMBA AXI 4 Stream" compliant back end interface wrapper
Synchronous clock design

Baseband PHYs

ntG3_BBP
Smart Grid PLC Baseband Processor
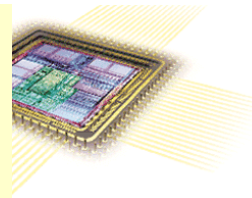
The ntG3_BBP is a fully compliant ITU-T G.9903 baseband modem that can be used in a wide range of smart grid applications over power lines, including smart metering and energy management in energy generation and distribution systems, lighting and industrial automation as well as automotive EV charging. The ntG3_BBP IP core main functional blocks are the transmitter, the receiver, the register file, the AHB-Lite wrapper and the analog front end interface. The user accesses the core via the AHB interface to either program the register file or provide payload data to the core. By programming the register file the user sets a specific functional mode of operation, requests the ntG3_BBP to transmit a data or acknowledgement PHY frame or accesses remotely received control information.
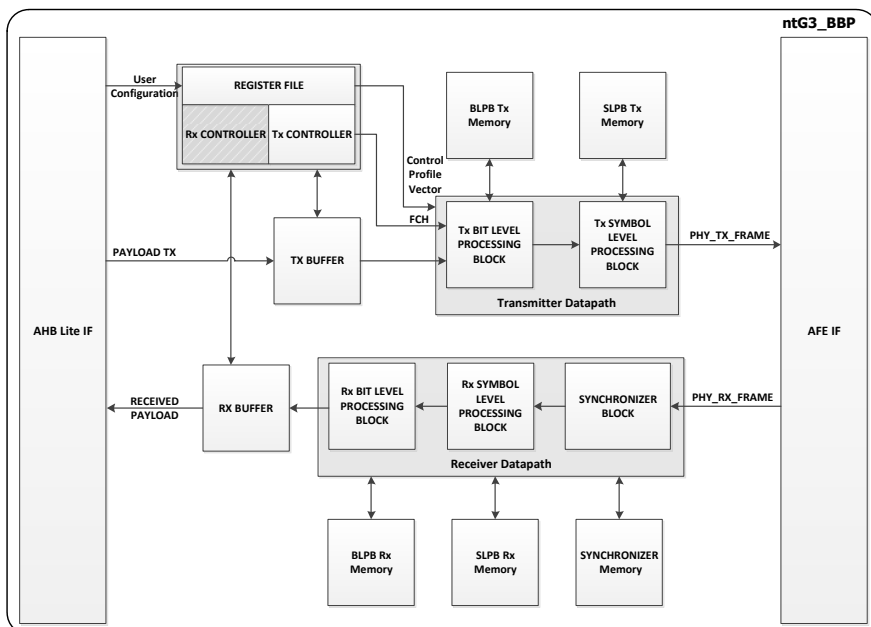
Features
PLC G3 physical layer (PHY) compliant baseband processor as per ITU-T G.9903 Chapter 7 and ITU-T G.9901 Annex B.
CENELEC-A/B (3-148.5kHz) and FCC (9-490kHz) bandplans support.
Aware of basic MAC layer handshaking primitives.
Data rates from a few kbps up to 290kbps.
On the fly programmable control profile selection.
Compliant with AMBA AHB-Lite protocol.
Synchronous clock design.
Silicon proven in Xilinx FPGA implementation technologies.

The Transmitter encodes the active register file control information to an FCH data block and computes all derivative control configuration parameters required for payload transmission. As per G.9903 both FCH and Payload data blocks are scrambled, Forward Error Correction (FEC) encoded, interleaved and finally modulated and tone mapped to OFDM symbols. These procedures form the bit level processing part of the Transmitter. After the tone mapper, the symbol level processing part of the transmission takes place. The OFDM symbols are optionally pre-emphasized, the IFFT transform is applied, and then the symbols are circular shifted and extended, windowed and overlapped to a concatenated final PHY Frame. The formed PHY Frame is provided to the AFE transmit interface, which synchronizes the symbols to the front end clock rate.

The AFE receive interface passes data from the channel to the Synchronizer unit. The Synchronizer is responsible of sensing the channel conditions and establishing a synchronization point in time, in other words it locates the start of the PHY Frame by searching for the Preamble symbols energy levels. After synchronization the Receiver removes the effects of windowing and the cyclic prefix, applies the FFT transform to cross the signal back to the frequency domain and thoroughly estimates the channel conditions in terms of phase and frequency offsets. Then the Receiver equalizes and compensates the received OFDM symbols accordingly.

Once the FCH OFDM part identified and equalized, it is being demodulated by the tone demapper and passed to the receive bit level processing block for de-interleaving, FEC decoding and de-scrambling. A CRC check is performed on the received FCH information and if successful, the Payload reception mode of operation is configured. All derivative control configuration parameters required for payload reception are computed. Then the Payload OFDM symbols are tone demapped, de-interleaved, FEC decoded and descrambled. Finally the decoded payload is being returned to the User via the AHB-Lite interface. The ntPLC_G3_BBP is able to implement basic coordination handshaking for interoperability with the MAC layer via designated fields of the Register File. TXOFF_RXON, BUSY_TX and BUSY_RX status flags are implemented.PD-DATA.request,PD-DATA.confirm, PD_ACK.request and PD_ACK.confirm status handshake primitives are supported as per G.9903 7.17.

The ntG3_BBP core has been synthesized using Xilinx ISE Design Suite tools. The core has been targeted to Kintex-7 XC7K410T-2 FFG900 device with a default balanced optimization strategy between area and timing. The implementation details are shown in the following table.

| Silicon Vendor	Device	Resources	Fmax (MHz)
Xilinx	Kintex 7 XC7K410T-2	22996 CLB Slices / 105 BRAMs / 87 DSP48s	81

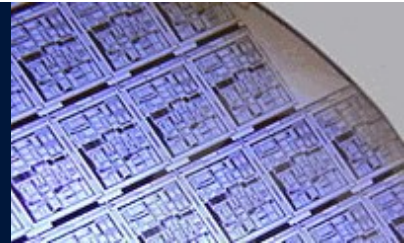# ntG3_BBP
Smart Grid PLC Baseband Processor
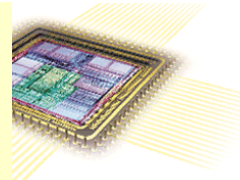
The ntG3_BBP is a highly programmable core, with numerous possible control configurations combinations of parameters such as : □ Band plan selection | □ Modulation Type and Mode | □ OFDM Symbols number | □ Legacy of Full Interleaver | □ 1 or 2 Reed Solomon codewords | □ Tone Map enabled and disabled sub-bands combinations | □ Tone Mask enabled and disabled tones combinations | □ Transmit gain control options
The ntG3_BBP IP Core achieves exceptional error correction performance as illustrated in the following FER vs Es/No graphs and is fully compliant with G3-PLC alliance performance masks (http://www.g3-plc.com/home/).
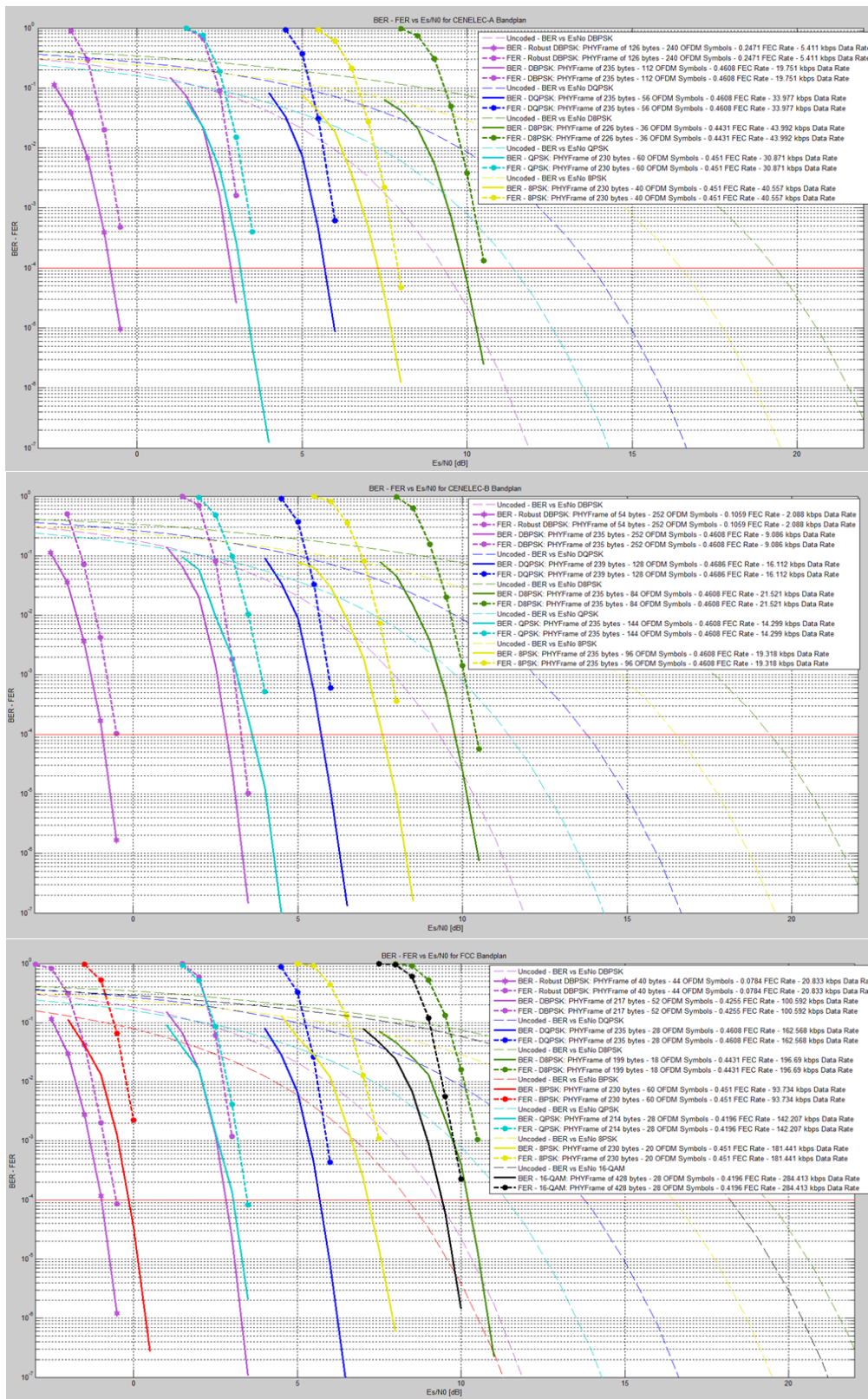
IP Customization—System Design—Consulting

Noesis Technologies offers expert ASIC, FPGA and DSP development resources to get your product in the market in time. Our highly skilled engineering team has considerable expertise in modeling, design and efficient implementation of telecom systems based on complex DSP algorithms. A list of our services include the following offerings:

► Telecom systems design feasibility analysis and specifications development.
► System level modeling in Matlab or System C.
► Efficient FPGA implementation and system demo prototyping.
► Customization of our existing IP Cores portfolio in terms of performance, interfacing and functionality.
► IP core development.
► Consulting on communications systems implementation in hardware.